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THESIS

A FRAMEWORK FOR FAULT TOLERANCE IN VIRTUALIZED SERVERS

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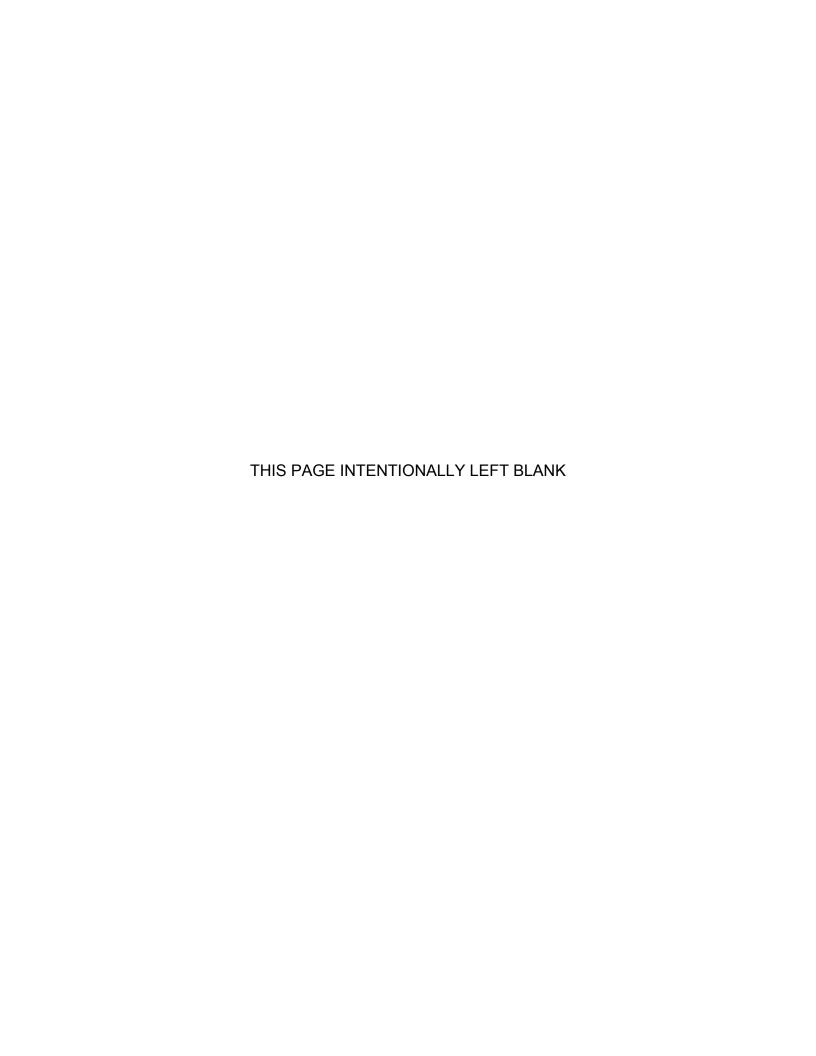
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June 2016

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A modern approach to ensuring the robustness of these systems is to have multiple systems (or servers) running at the same time that back up one another. Since that approach is expensive, this thesis attempts to solve that problem or find an alternative solution with a fault-tolerant, virtual server-based system framework. Our goal is to overcome shortcomings with a cost- and space-efficient and user-friendly approach.

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A FRAMEWORK FOR FAULT TOLERANCE IN VIRTUALIZED SERVERS

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ABSTRACT

In modern naval platforms, most of the critical operations are done with the help of automated systems. Specifically, operational decisions and actions are finalized using command and control systems (C2 systems). A wide variety of sensors, radars, communication devices, and weapons are connected to C2 systems. Generally speaking, C2 systems receive data from their respective sensors and radar and process that data. Officers then rely on C2 output to make sound decisions by using their technical knowledge combined with detailed scientific information.

A modern approach to ensuring the robustness of these systems is to have multiple systems (or servers) running at the same time that back up one another. Since that approach is expensive, this thesis attempts to solve that problem by finding an alternative solution with a fault-tolerant, virtual server-based system framework. Our goal is to overcome shortcomings with a cost- and space-efficient and user-friendly approach.

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LIST OF ACRONYMS AND ABBREVIATIONS

C2 command and control
DBA database administrator

DGMGRL Data Guard command line interface

HDFS Hadoop Distributed File System

RAID Redundant Array of Inexpensive Disks

URL Uniform Resource Identifier

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I. INTRODUCTION

This chapter is an introduction to the thesis. The background in the research area, the purpose of the thesis, its scope and limitations, thesis question, and the organization of the thesis are presented in this chapter.

A. BACKGROUND

Navy platforms operate under all weather and sea conditions. Therefore, it is imperative that the systems running on such platforms be robust. In the real world, however, this is hard to achieve since there are many radars and sensors connected to the majority of systems. Moreover, since many users operate a single system simultaneously and in shifts, this burden exacerbates the potential for error. Yet even under these circumstances, active Navy platforms' systems have to be resilient and reliable—up almost all the time. That means the systems have to be fault-tolerant against users, errors from sensors/radars, system errors, and physical or hardware errors.

As technology has developed during the past four decades, most of the equipment used on Navy platforms has become more dependent on computers. In today's Navy platforms, most of the critical operations are run with the help of automated systems. The designed automated systems have to be robust, fault tolerant, and have high survivability. Otherwise, rather than being freed up to focus on operations, personnel will be beleaguered with the problems that occur in the automated system.

A modern approach to ensuring the robustness of these systems is to have multiple systems (or servers) running at the same time that back up one another. This thesis attempts to solve these problems or find an alternative solution with a fault-tolerant, virtual server-based system framework. Our goal is to overcome shortcomings with a cost and space-efficient, user-friendly approach.

B. PURPOSE OF THIS THESIS

The main purpose of this thesis is to find an alternative way to ensure the fault tolerance of data storage and retrieval in command and control systems (C2 systems) of Navy platforms, especially when the data is in transit. By doing this, when a problem occurs in system data storage and distribution units, the C2 system will find an alternative way to stabilize itself and recover from an unintended state to a normal one. The concept of fault tolerance has to be implemented in the C2 system according to a plan to ensure the system will behave as expected in wartime under any circumstances. Thus, an implementation plan must include these essential considerations:

- The C2 system itself must have no single point of failure. This also means that there must always be a way for the system to compensate for failures.
- Although it is impossible to have a zero-failure environment, the C2 system must maintain stability and have a flawless runtime by being able to rapidly switch to alternative data resources at the time of failure.
- The required time to switch between main data resource and alternative should be minimal.
- Switching between data resources at the time of failure must not affect the rest of the system, especially ongoing processes of high importance.
- The C2 system needs to recover itself instantly. Moreover, it also needs to check the status of the main data resource frequently to switch back to that resource as soon as the problem is resolved.

This thesis tries to answer the following question:

 Is there a better way to ensure the fault tolerance in command and control systems of Navy platform databases in terms of economics and performance?

C. SCOPE AND LIMITATIONS OF THIS THESIS

The scope of this thesis is to find an alternative way to have a fault tolerance in databases in C2 systems of Navy platforms using Oracle Database 12c and its features and to test it respectively according to fault tolerance

concepts. This thesis mainly focuses on Oracle Database 12c and Oracle Data Guard, which run on the Fedora operating system.

In this thesis, we used Oracle VirtualBox as a hypervisor for the virtual environment that we worked on. We also used Fedora operating system to install the Oracle Database 12c. We created two databases in the system. Those two databases use the localhost as host and different service names to communicate via Data Guard. They also have four standby redo log files, which have the size of 50MB each. The listener and Data Guard command-line interfaces (DGMGRL) are used for the communication of the system. The observer is installed on the same operating system.

In addition to those system characteristics previously mentioned, we also connected to the databases as a normal user (in other words, as a client) using Java codes on the same operating system. In this framework, instead of making the client failover scenarios transparent to clients using one single connection URL (Uniform Resource Identifier), we explicitly used two URL strings in order to connect to the databases so that we can measure the time to wait for changing the connection URL in between failover or switchover testing scenarios by extensive tests that are independently calculated in the system.

D. ORGANIZATION OF THIS THESIS

Chapter I provides a general overview on the thesis. In Chapter II, we cover the fault tolerance concept by introducing definitions of fault, system failure, and error as well as many other definitions related to the fault tolerance concept. After presenting those definitions, we review some technologies that ensure fault tolerance. After that, we introduce Error Detection and Error Correction subjects as the stages of fault tolerance. We conclude this chapter with five major objectives of fault tolerance: Reliability, Availability, Safety, Maintainability, and Dependability.

In Chapter III, we specifically focus on the fault tolerance in Oracle Databases. We explain the Data Guard feature in this chapter. We also present how Data Guard synchronizes databases in this part of the thesis.

In Chapter IV, we introduce the design and implementation of the system for fault tolerance in databases. After introducing the design and implementation, we present various testing cases and findings.

In Chapter V, we conclude the study and recommend some ideas for future studies.

II. A REVIEW OF FAULT TOLERANCE

In this chapter, definitions of fault, system failure, and error as well as many other definitions related to the fault tolerance concept are discussed. After presenting those definitions, we review some technologies that ensure fault tolerance. Hot Swapping, RAID (Redundant Array of Inexpensive Disks), and Apache Hadoop technologies are presented under the technologies section. After that, Error Detection and Error Correction subjects are introduced as stages of fault tolerance. This chapter concludes with five major objectives of fault tolerance: Reliability, Availability, Safety, Maintainability, and Dependability.

A. AN OVERVIEW OF FAULT CONCEPT

In this part of the thesis, introducing the breakages of dependable computing in the first place is a good starting point. There are a few impairments to dependable computing such as faults, errors, and failures in the system [1].

A user can only perceive the part of the system that is being interacted with as the system behavior. A normal system behavior would be getting the expected service back from the system. For example, when clicking on the "Save" button on a Microsoft Word document, if the system is saving the document that the user wanted to save, then that is the normal system behavior. Another example, when the user wants to make a simple calculation using a calculator, the normal system behavior should be getting the correct outcome after the user presses the "=" button.

On the other hand, a system can respond with an unusual or unintended behavior. This can be explained further using the examples just given, but ending with different results. When clicking on the "Save" button on a Microsoft Word document, the system does not save the document that the user wanted saved, or the calculator produces the wrong outcome after the user presses the "=" button.

The service that is delivered by a system can also be interchangeably called the system behavior. A system failure happens when the delivered service is not the same as the specified service, as opposed to the normal system behavior. A system failure is a specific type of unusual behavior, because it is not intended. An error is a part of the system itself, and an error may also be the underlying cause of the failure. An error, which is introduced into the system by a programmer, may lead to a fault, and that fault can cause a system failure [1], [2].

Here are some related definitions taken from the *IEEE Standard Glossary* of *Software Engineering Terminology*:

- Error: The difference between a computed, observed, or measured value or condition and the true, specified, or theoretically correct value or condition. For example, a difference of 30 meters between a computed result and the correct result.
- Fault: An incorrect step, process, or data definition. For example, an incorrect instruction in a computer program.
- Failure: An incorrect result. For example, a computed result of 12 when the correct result is 10.[2]

To put the definitions in an order for a better understanding, here is the typical system behavior when there is a system failure: an error creates a hidden fault, which becomes effective after its activation, and the system failure happens when the fault, which was hidden beforehand, affects the normal system behavior or the service. In other words, a fault is the result or the appearance of an error in the system, and a failure is the result or the appearance of a fault in the system [1], [2].

A fault, in a general sense, is a kind of physical defect that happens in the software or hardware parts of the system [3]. Since a system is too big and complex, it is impossible to test every single sub-system and part—especially when no failure is observed. For this reason, it is quite hard to foresee a fault. When there is a fault in the system, it is not certain that there is a failure in the system. A failure only happens if a hidden error is activated and causes a fault; then a hidden fault is activated and causes a failure to happen. On the other

hand, a failure in the system assures that there is a fault. Therefore, an error is present in the system.

Here are some more examples on the notion of error and fault concept:

- A programmer's mistake, which is in the code, is a kind of error.
 The consequence of that error is a hidden fault that is in the
 program. That particular fault stays hidden or silent until it is
 activated. At that point, that fault creates an erroneous data, thus, a
 failure happens [1].
- A hardware defect in hard disk is a kind of fault. That particular fault stays hidden or silent until it is activated. At that point, that fault creates an erroneous data, thus, a failure happens.
- Like the previous example, there are some other similar types of fault such as an electromagnetic unsteadiness because of change of energy and an inappropriate man-machine interaction [1].
- A mistake, typographical error, or misleading instructions in a maintenance or operating manual is a kind of error. Actually, this example is a very good one, since it shows us that failures are not always about software or at hardware level. Failures can be related to operators, weather conditions, and even documentation, in addition to software and hardware infrastructure. It is simply because we cannot separate any system from the outer environment, or from people who are interacting with them. In this maintenance or operating manual example, the error will stay as hidden in the maintenance or operating manual until the directive, which has an error in it, is applied to the system [1], [2].

There are various reasons behind faults. Examples include:

- Physical faults: short-circuits in the hardware level, temperature, vibration, heat, humidity, etc[1]..
- Human-made faults: generally human-made faults happen from some other reason, such as the following:
- Design faults: this type of fault is generally made during system design or system modifications, or during the creation of maintenance or operating manuals.
- Interaction faults: deliberate or non-deliberate faults of operators.
 This type of error may occur because of the lack of personnel training [1].

B. AN OVERVIEW OF THE FAULT TOLERANCE CONCEPT

In this section, the concept of fault tolerance is presented. It is a highly important subject in the modern computing world, because there are numerous activities, including those that shape countries' economies and—even more important—those that protect human lives, that should be running with minimal, or if possible, no faults.

1. Definitions

A system is fault tolerant if it can tolerate or minimize the effects of the existing faults in the system by the help of redundancy [4]. Actually, it is easier said than done. The system cannot be made fault tolerant against every single fault as a whole. The utmost aim of fault tolerance is avoiding the total system failure when some sub-parts or sub-systems fail [4].

Another definition of fault tolerance can be given as a system is treated as fault tolerant provided that the system behavior or, in other words, expected service is consistent with the normal system behavior, even when there are some failures in the system [4].

Having a fault tolerance in the system may have some minor adverse effects such as decline in the overall performance or limitations of total disk capacity. However, these effects are not so important in comparison to the failure of the total system.

Redundancy is the key behind the fault tolerance concept [4]. From another perspective, fault tolerance cannot be offered without any means of redundancy [4]. Redundancy is defined as a kind of abundance of resources, where those resources would be not used or not needed as the system is running without faults [4]. More about the redundancy concept is presented in the next section.

There are a few more definitions related to the fault tolerance discussion worth noting:

- Fault avoidance is preventing the possibility of fault occurrence.
- Error removal is minimizing the existence of hidden errors in the system by using verification.
- Error forecasting is estimating the presence, occurrence and consequences of errors in the system [1].

Fault tolerance and fault avoidance are mainly part of dependability acquisition, which means giving the best effort to deliver the specified service. On the other hand, error removal and error forecasting are mainly part of dependability validation, which means getting enough confidence off the system on delivering the specified service [1].

2. Redundancy Concept

Under normal operation, a redundant part of a system is abundant and does not participate with the parts that are actively involved in the system run [4], [5].

According to Koren, "Redundancy is the property of having more of a resource than is minimally necessary to do the job at hand. As failures happen, redundancy is exploited to mask or otherwise work around these failures, thus maintaining the desired level of functionality." [6]

To qualify as a fault tolerant system, the system has to have redundancy. Redundant parts are to be used in place of a failed part when there is a failure in the system. Redundancy may bring some side effects into the system. Decrease in performance, the expansion in the total system size and weight, and a hike in the system cost can be counted in those adverse effects [5].

There are four types of redundancy in a general aspect. These redundancies are hardware redundancy, software redundancy, time redundancy, and information redundancy [5], [6].

Hardware redundancy can be established by having extra or, in other words, abundant hardware in the system for mainly two purposes. One is to detect the effects of a failed component in the system, as in run-time testing

uses. The other is overriding the effects of a failed part. For example, we can have two or more power supplies, each performing the same function by giving adequate redundancy to ensure our system will run as intended, instead of having a single power supply in our system. We can detect the error in a running power supply if we have two power supplies. If we have three power supplies, we can override the effects of a failed power supply, and so on. We can give the power supply instance just presented as a static hardware redundancy example. The main goal of static hardware redundancy is to mask the failure immediately by allowing the system to run as if there is no failure at all. There is another type of hardware redundancy, which is called dynamic redundancy, which activates the spare components in the system as a failure occurs. In this dynamic hardware redundancy, spare components in the system are not used during normal operation, and they will be waiting passively until there is a failure and the need for them occurs. Both static and dynamic redundancy approaches for fulfilling the users and systems needs are called hybrid hardware redundancy [6].

3. Some Technologies that Ensure Fault Tolerance

Some technologies have been developed over the years to ensure the fault tolerance in a given system. In this section, Hot Swapping, RAID, and Apache Hadoop system is presented here.

a. Hot Swapping

Hot swapping is a general term to describe the functions that change some parts without shutting down the whole system. Hot swapping function also keeps the system running with no interruption or only a minor interruption at the worst case. There are two main types of hot swapping, known as hardware hot swapping and software hot swapping [7], [8].

The idea behind hardware hot swapping is to have abundant disks or hardware components in order to keep the system running normally, even if there is a failure. Hot swapping is strictly a hardware function and requires no commands from the terminal or manual operation. If this were the case of

applying a manual procedure or putting the system in sleep mode or in any state other than the normal one, then this would not be a fault tolerant system at all.

According to [8], software hot swapping is the technology term that explains the "replacement of a software program or a part of a program while the whole software system remains in operation" [8]. Hot swapping indicates the ability to replace the parts that are not functioning without the user's notice [7].

According to Feng [8], "Software upgrading, needed for bug-fixes, updates, or functionality upgrades, is generally not easy. It is particularly difficult in computer communication networks where software can be widely distributed across heterogeneous domains. The problem is complicated even further by the requirement in some applications for almost 100% availability. As a result, in many cases it is extremely important not to have to take a system off line for software upgrading and/or recompilation."

There are some issues in software hot swapping, such as the Referential Transparency Problem, the State Transfer Problem, and the Mutual Referential Problem [8], [9].

b. RAID

RAID is the abbreviation for "Redundant Array of Inexpensive Disks." It is a technology for storing data that associates multiple physical hard drives into a logically, or seemingly, single unit. This technology is designed primarily to provide fault tolerance, large storage capacity, and faster disk access. The main focus areas in this technology are performance and reliability [10].

There were five standard RAID levels originally; however, many more have evolved as time passed [10]. The standard RAID levels are named as RAID and a number starting from 0. For example, standard RAID levels are named as RAID 0, RAID 1, and so on. In addition to standard RAID levels, it is possible to create hybrid levels. The most common examples of hybrid RAID levels are just one level deep. In other words, generally two standard levels are combined to

create hybrid RAID levels. Combined RAID levels are named as RAID and the combination of numbers that forms the specific RAID combination. For example, combined RAID levels are named as RAID 0+1, or RAID 01, if that combined RAID level is the combination of RAID 0 and RAID 1 [10].

Standard RAID levels 0 through 6 are illustrated in Figure 1.

RAID Levels 0 through 6 Figure 1. Non-Redundant (RAID Level 0) Mirrored (RAID Level 1) Memory-Style ECC (RAID Level 2) Bit-Interleaved Parity (RAID Level 3) Block-Interleaved Parity (RAID Level 4) Block-Interleaved Distributed-Parity (RAID Level 5)

P+Q Redundancy (RAID Level 6)

Source: P. Chen P et al., "RAID: High-performance, reliable secondary storage," *ACM Computing Surveys (CSUR)* vol. 26, no. 2, p. 153, 1994.

RAID 0 is disk striping. In RAID 0, there are multiple physical hard disks that form a big single disk. For example, if we have four 1 GB hard disks in our system and we use RAID 0, then we have 4 GB of total capacity. In other words, we sum the total capacities of each disk in order to find the capacity of our RAID 0 volume. RAID 0 does not offer any type of redundancy for disk failures. If a disk failure happens, then all the data in that disk is practically lost.

Disk mirroring technology is known as RAID 1. In RAID 1, in short, multiple physical disks look like one disk. This technology will duplicate the same data onto a paired disk; yet, they will look like one disk. For a simple example, assuming that we have four 1 GB hard disks in our system and we use RAID 1, then we could have two pairs of disks for creating real-time disk copy; therefore, the whole disk capacity may look like 2 GB instead of 1 GB. We could visualize the same example as we do have two pairs, each having two 1 GB disks in it, and each pair looks like one disk. RAID 1 provides fault tolerance by redundancy. If one disk fails, we can still work because the other disk in the pair (by saying 'pair' here, we are trying to refer to the example that we gave earlier) has the exact copy of the failed disk. We can also change the damaged or broken hard disk while the system is running. RAID 1 enables the hot swapping technology, which is presented in the previous subsection. The biggest disadvantage of RAID 1 is that we lose half of our performance and half of our total disk capacity. RAID 1 technology does not hold a backup for users. It only works on hard disk failures, where it offers fault tolerance by redundancy. For example, when a user deletes a file from the hard disk, the exact copy is deleted from the disk next to it (that is, from the other disk holding the exact copy of the previous disk).

RAID 5 is known as "block-level striping with distributed parity" [11]. This RAID level has a place in between RAID 0 and RAID 1. We can reconstruct the data after failing, with the help of parity bits and error correction codes. If a disk fails, we can replace it before any other disks fail and continue our work, and the user can feel the degraded performance. According to [11], "Upon failure of a

single drive, subsequent reads can be calculated from the distributed parity such that no data is lost. RAID 5 requires at least three disks."

c. Apache Hadoop

Apache Hadoop is a "framework that allows for the distributed processing of large data sets across clusters of computers using simple programming models" [12].

The core of Apache Hadoop has two main parts as the storage part and the processing part. The storage part is called the Hadoop Distributed File System. The processing part is called the MapReduce.

Hadoop Distributed File System (HDFS) stores large files that are generally in the range of gigabytes to terabytes across multiple machines [13]. HDFS is managed by daemon processes, which are as follows:

NameNode: Master process

DataNode: Slave process

Checkpoint NameNode (or Secondary NameNode): Checkpoint process

BackupNode: Backup NameNode [13]

HDFS architecture is presented in Figure 2.

Client

3. Read/write data

DataNodes

Figure 2. HDFS Architecture

Source: D. Borthakur. (2011, Apr. 04). Apache Hadoop filesystem and its usage in Facebook. [Online]. Available: http://cloud.berkeley.edu/data/hdfs.pdf

NameNode is the master process daemon server in HDFS that coordinates all the operations related to storage in Hadoop, including the reads and writes in HDFS. NameNode manages the filesystem namespace. NameNode holds the metadata about all the file blocks, and in which all nodes of data blocks are present in the cluster [13].

DataNode holds the actual data in HDFS and is also responsible for creating, deleting, and replicating data blocks, as assigned by NameNode [13].

Checkpoint NameNode, earlier known as Secondary NameNode, is a node that has frequent data check points of FsImage and EditLog files merged and available for NameNode in case of any NameNode failure. Checkpoint NameNode collects and stores the latest checkpoint [13].

BackupNode is similar to Checkpoint NameNode, but it keeps the updated copy of FsImage in RAM memory and is always synchronized with NameNode [13].

MapReduce is a

programming model and an associated implementation for processing and generating large data sets. Users specify a map function that processes a key/value pair to generate a set of intermediate key/value pairs, and a reduce function that merges all intermediate values associated with the same intermediate key [14].

MapReduce provides "automatic parallelization and distribution," "fault tolerance," "input and output scheduling," and "status monitoring" [14].

MapReduce, which is fault tolerant software, is designed to avoid server problems. According to Dean, "When a machine fails, the master knows what task that machine was assigned and will direct the other machines to take up the map task. You can end up losing 100 map tasks, but can have 100 machines pick up those tasks" [15].

According to Shankland,

The MapReduce reliability was severely tested once during a maintenance operation on one cluster with 1,800 servers. Workers unplugged groups of 80 machines at a time, during which the other 1,720 machines would pick up the slack. "It ran a little slowly, but it all completed," Dean said. ... In a 2004 presentation, Dean said, "One system withstood a failure of 1,600 servers in a 1,800-unit cluster" [15].

C. OBJECTIVES OF FAULT TOLERANCE

Fault tolerance is an attribute built in a system that ultimately seeks to meet design requirements. The most significant requirements are dependability, availability, reliability, safety, and maintainability [3], [5].

1. Dependability

The definition of dependability according to Avizienis et al. is "the ability to deliver service that can justifiably be trusted. This definition stresses the need for justification of trust. The alternate definition that provides the criterion for deciding if the service is dependable is the dependability of a system is the ability to avoid service failures that are more frequent and more severe than is acceptable" [16].

Avizienis et al. also state, "It is usual to say that the dependability of a system should suffice for the dependence being placed on that system. The dependence of system A on system B, thus, represents the extent to which system A's dependability is (or would be) affected by that of System B. The concept of dependence leads to that of trust, which can very conveniently be defined as accepted dependence" [16].

Dependability is a concept that integrates and covers some other concepts in fault tolerance, which are summarized by [16] as:

- Availability: the state of being available for desired service
- Reliability: ensuring that the ongoing service is correct
- Safety: ensuring that no catastrophic consequences will result to users and the system
- Maintainability: the flexibility of going into small or big changes, yet still being able to function

2. Availability

The availability of a system can be expressed with a function, expressed as the probability that "the system is operating correctly and available" to its users at time t [3]. Availability is an important metric for systems. Most of the big technology enterprises introduce their products with an availability metric, as well as the specifications of that product. Figure 3 gives the availability chart for systems in general.

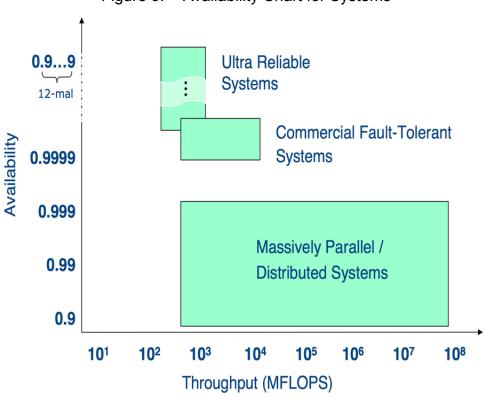


Figure 3. Availability Chart for Systems

Source: M. Malek. (2004, May). Dependable systems introduction. [Online]. Available: http://www2.informatik.hu-berlin.de/rok/zs/WS0405/data/slides/zs01_04Intro.pdf

System availability can be given in a chart after calculating the ratio of the time where the system is available to the total. Figure 4 shows some common classes with the associated availability percentages and annual downtime of the system [5]. Systems are named after the total number of 9's they have in their availability measurement. For example, if a system has the value of 99.9% for system availability, then the system can be named the three nines system. Alternatively, the system can be expressed as an availability class. For instance, if a system has a value of 99.995% for system availability, then the system can be categorized as Class 4, since 99.995% is between 99.99% and 99.999% [5].

Figure 4. Availability Classes

AVAILABILITY MEASUREMENT	ANNUAL OUTAGE	AVAILABILITY CLASS	Č.
90%	More than a month	One nines	G) 1
99%	Just under four days	Two nines	Class 1 Class 2 Class 3 Class 4 Class 5 Class 6
99.9%	Just under nine hours	Three nines	
99.99%	About an hour	Four nines	
99.999%	A little over five minutes	Five nines	
99.9999%	About half a minute	Six nines	
99.99999%	About three seconds	Seven nines	

Source: E. Kati, "Fault-tolerant approach for deploying server agent-based active network management (SAAM) server in Windows NT environment to provide uninterrupted services to routers in case of server failure(s)," M.S. thesis, Dept. Computer Science, Naval Postgraduate School, Monterey, CA, 2000.

3. Reliability

The reliability of a system is a "function of time, R(t), defined as the conditional probability that the system performs correctly throughout the interval of time, $[t_0,t]$, given that the system was performing correctly at a time t" [3]. In other words, it is a measure of the "continuous service accomplishment from a reference initial instant" [3].

According to Sorin [17], "Unless a system failure is catastrophic (e.g., avionics), reliability is a less useful metric than availability."

The reliability of a system can be expressed with the 'Bathtub Curve,' which is given in Figure 5. Generally, there are three phases in a system's life in terms of failure rate, which are decreasing failure rate, constant failure rate, and increasing failure rate. It is quite important to determine the time projections,

where the decreasing failure rate curve ends and the constant failure rate starts, as well as where the constant failure rate ends and increasing failure rate curve starts. This has two important advantages for the system. One is the advantage of determining the time projections correctly, which enables the manufacturer to decide the length of the guarantee time or the warranty. Most of the guarantee time lengths are determined according to failure rates of the system. In the user's perspective, it is always safe to have a warranty or extended warranty right before the failure rate increases.

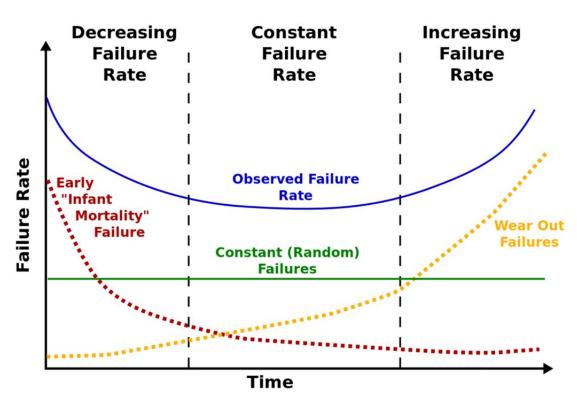


Figure 5. Bathtub Curve for Hazard Function

Source: Bathtub curve. (n.d.). *Wikipedia.* [Online]. Available: https://en.wikipedia.org/wiki/Bathtub_curve. Accessed Dec. 24, 2015.

The other advantage is determining the time projections correctly, which gives the knowledge of when to replace the subparts of a system or completely

remove and replace with the upgraded version of the system before the failure happens, especially in the military domain, where almost all systems are critical.

Assuming that the system life time is exponentially distributed, the reliability of that system is $e^{-\lambda t}$ [5]. In the formula, λ is the failure rate of the system and t is the time interval. The expected failure numbers per the unit of time is called the failure rate of that system [5].

According to Pradhan, "The exponential relationship between the reliability and the time is known as the exponential failure law. The exponential failure law is very important for the analysis of electronic components, and is by far the most commonly used relationship between reliability and time" [3], [5].

4. Safety

According to Pradhan, "Safety, S(t), is the probability that a system will either perform its functions correctly or will discontinue its functions in a manner that does not disrupt the other systems or compromise the safety of any people associated with the system" [3]. Safety is a representation of the fail-safe capability in a system. If the system does not function as expected, we expect it to fail in the safest way [3], [5].

The concept of reliability differs from the concept of safety. Reliability can be expressed as a probability of any system's correct overall function performance. On the other hand, safety tells that the system will either function correctly or will run without interrupting the whole process. The concept of safety is expressed with a probability. It is astute to say that if a system is reliable, it is also safe. On the other hand, we cannot conclude the other way around [5].

5. Maintainability

According to Pradhan [3], "The maintainability is a measure of the ease with which a system can be repaired once it has failed." It is the "probability that a failed system will be restored to an operational state" within a certain and acceptable period of time [3], [5].

According to Kati, "Maintainability encapsulates not only the failures of the system, but also the modifications that are necessary for the required level of system performance" [5]. System functions must be maintained and upgraded regularly to ensure that the system meets user expectations and needs. Performing these maintenance activities can be made easier if the system is highly modular [5]. Kati further states that "If the consequences of a modification can be localized to well-defined small modules, then the maintenance effort can be minimized. Fault tolerance can support maintainability in the problem detection and problem location process. The maintenance of the system can be done after the detection and the localization of the failure. Fault tolerance can also support maintainability in the modification process by allowing maintenance actions without interrupting the service delivered by the system" [5].

D. FAULT TOLERANCE STAGES

The systems that have fault tolerance try to continue providing normal service despite having failures in some subcomponents. The most important parts of fault tolerance are error detection, damage confinement, error recovery, and fault treatment and continued system service [18].

Error detection is the first stage, where a fault and, therefore, an error is detected in a subcomponent. The detection of an existing error implies that a failure may occur on that subcomponent, as discussed in the section "An Overview of the Fault Concept," earlier. After detecting the error, the system has to identify and bound the limits of the damage that was the consequence of the failure. This stage is known as damage confinement. In these first two stages, the error is detected, and efforts for limiting the estimated boundaries of the damage caused by the failure are endeavored. These first two stages can be grouped as detection stages.

Errors and faults have to be detected before starting the further work for fault tolerance. When the system finishes these first two stages, errors and faults need to be corrected for enabling the normal service to be delivered to users.

This is the responsibility of the error recovery stage. Error recovery stage will take the system to an error-free state, by removing errors that occurred beforehand.

After the error recovery stage, the fault treatment and continued system service stage comes into action for identifying the faulty component or components.

1. Error Detection

In the extent of fault tolerance, error detection is the very first step. In the best case, we expect the error detection mechanism to detect every single error in the system, which is only ideally true. In reality, however, it is not quite possible.

There are certain features that an ideally feasible error detection mechanism should fulfill. First, the error detection should not be affected by the system design. Therefore, the best approach is seeing the system as a black box, like a function structure in any programming language [4], [5].

Second, the error detection mechanism should be sound and complete. This can be achieved by detecting all possible errors, and the declared errors should indeed be in the system [4], [5]. The detection mechanism should avoid giving false positives [4], [5].

Third, the check should be independent from the system in terms of being vulnerable to failures. The detection mechanism should not fail at system failures; otherwise, the whole error detection phase would be unproductive and meaningless, which amounts to having a different error detection failure mechanism than the actual system [4].

These three criteria can rarely be met in systems in real life due to many difficulties. It is very hard to identify every single error in the system. The complete test of error detection is very costly and time-consuming [4].

Because of the aforementioned reasons, in most of systems, checks for acceptability are made instead of checks for ideality. In other words, checks for acceptability are the approximation of checks for ideality. The main aim in this perspective is to lower the cost of error detection checks, as well as trying to keep the error detection performance maximized. In the checks for acceptability approach, most errors in the system are expected to be caught in a predefined confidence interval. Here are some general types of error detection checks that are used frequently:

(1) Replication Checks

Replication checks are one of the most broadly used and effective checks. These types of checks have advantages in terms of ease of use and completeness in checks [4]. They can also be applied to a system with a little knowledge of internal system dynamics. As it may be understood from the name, these checks depend on replicating some components of the system. Therefore, replication checks are highly expensive error detection methods [4]. Replication, which is using identical copies of system parts, works provided that the interior design of those parts is working as expected. On the other hand, if the interior design of those parts is not working normally, replication checks will not succeed either [4].

Replication has purposes other than solely for error detection, especially in distributed systems [19]. Generally, data or processes, which are replicated in distributed systems, enable the handling failures in the system easily [4].

(2) Timing Checks

If timing constraints are included in the interior design of system components, timing checks can be used for checking whether timing constraints are satisfied or not. Timing checks mainly play a big role both in hardware and software systems for detecting problems. After setting the timer, the system will check whether the timer has timed out. Having a timeout in the system indicates

that there is at least one component malfunctioning. In other words, a timing error is an indirect indication of the existence of an error in the system [4].

(3) Structural and Coding Checks

Two types of checks are used in general, semantic checks and structural checks. Semantic checks are used for verifying the value is consistent with the rest of the system. Structural checks, on the other hand, are used for checking data against data, which is the part of the normal behavior of the system, in the internal design of system structure [4].

Structural checks are mainly used in hardware components of systems by means of coding. Extra bits are added to actual data bits for error checking purposes. When an inconsistency is found in coding bits with the help of extra bits, the error is detected. These kinds of error detection mechanisms are mainly used in hardware components of systems. Data structures that use redundancy for enabling structural checks are known as robust data structures [20]. In robust data structures, the system, which uses redundancy and structural checks, can locate the error and also correct the corrupted part by using extra bits [4]. This approach is used in RAID technology.

(4) Reasonableness Checks

Reasonableness checks decide whether the system status is reasonable. For example, the range check, which is made to check whether the value is in a specified range, is a type of reasonableness check. The range check only shows whether the value is within the range. It does not say anything about the correctness of the given value [4].

Inserting assertion statements in the code is another kind of reasonableness check. By definition, an assertion is a logical expression on the value of different variables in the system, which will evaluate to true if the state of the system is consistent; otherwise, it will evaluate to false. It is a practical way to detect errors, especially in software [4].

(5) Diagnostic Checks

The purpose of diagnostic checks is to determine whether the component of a system is working correctly. In these types of checks, the system performs diagnostic tests on a specific component. Systems mostly use previous known correct values to check current values for comparison. Therefore, a mismatch in comparisons equates to an error. Diagnostic checks, which are also known as self-checks, mostly appear on startup times of systems [4].

2. Damage Confinement

When an error in the system is detected, it also points out one or more faults that are present in the system. While it is good to detect a fault or faults in the system, we need to deal with the problem of time delay as well. The time delay in this context can be addressed as the time difference between the failure and the time that we detected the error. One of the many reasons for the time delay issue can be that we do not monitor the system all the time. Having some problems in our log files can be the other reason. One way or another, the time delay issue can lead to other problems in the other parts of system [4].

In this damage confinement phase, the boundaries of corruption have to be determined before going ahead and correcting actual failures.

3. Error Recovery

After detecting the error in the system and determining the boundaries of corruption, it is appropriate to start removing the error from the system. The main goal of this phase is to clean the system of errors. It is very important to apply error recovery on the system, so the system can lead to a consistent state. Two techniques for error recovery are backward and forward recovery [4].

(1) Backward Recovery

When we exercise backward recovery, the system is restored to an earlier state, where there is no error. In order to apply this error recovery technique, the

states of the system have to be saved periodically [4]. It is like taking backups, such as full, incremental, or differential backups, of the system. We can call those backup media checkpoints of the system. After isolating errors in the system, we can roll back to a previous backup, ideally the most current one that is known to be good. It is worth noting that the system has to be backed up frequently to allow for successful backward recovery in the long term.

The biggest advantage of this approach is that we do not need to analyze the possible causes of errors in order to apply the error recovery on the system. Thus, error recovery will not take too much time. We will only roll our system back to an earlier backup. It is also possible that we can investigate the reason for the errors on a different test machine, while our actual system is running normally. In large enterprises, this error-analyzing job can be given to special departments.

(2) Forward Recovery

Unlike backward recovery, there is obviously no backup available in the forward recovery technique. Therefore, this approach attempts to neutralize errors in the system and make it run normally. We achieve this by correcting existing errors [4].

It is quite obvious that the cause of an error has to be diagnosed and learned before moving ahead and starting forward recovery. Consequently, a very detailed error diagnosis has to be performed for forward recovery.

A special team of people has to work on error diagnosis for forward recovery.

4. Fault Treatment and Continued System Service

The main focus in the first three phases is detecting errors, determining the boundaries of corruption, and then removing the error from the system. After the first three phases, the system should be free from errors. If the error is temporary, or if it originated from a momentary variation in current, voltage, or frequency, error detection, damage confinement, and error recovery will be enough for the system to go back to an error-free state. Moreover, we can be sure that when we reboot the system, it will run normally if temporary or transient error is the case [4].

On the other hand, there may be some cases where the existing error is the result of some permanent condition. If that is the case, even if we remove the error from the system, after the restart or some amount of time; the same error will reoccur. Therefore, we need to identify the faulty component and take that component out of use in order to avoid the same error recurring repeatedly. The main goal of this phase is to bypass the faulty component of the system without affecting the normal runtime of the environment [4].

This fault tolerance stage has two main subparts, fault location and system repair.

(1) Fault Location

In the fault location subpart, the faulty component has to be identified. Unless the faulty component is found, full recovery cannot be made [4].

(2) System Repair

In the system repair subpart, the system will be repaired by either replacing the faulty component with a new one or taking the faulty component out of the system. In order to have a fault tolerant system, the system has to be online during the process of system repair. If we must power off the system during the system repair, we cannot call that system fault tolerant. Thus, the system repair has to be done in an automated manner. The redundancy concept plays a big role in this subpart [4].

One approach is having a working, errorless spare component in the system so that the system will have a spare component that has the exact same settings and properties as the affected component when a faulty component has to be replaced [4].

Another approach is virtualizing the whole system. This approach is cheaper than having a spare component in the system all the time. For some projects, despite the cost, it is preferable to have an abundance of software components for providing better system up time and more speed in terms of overall system performance.

After we apply the system repair technique, the system will run normally. The drawback of system repair may be a certain level of performance degradation. On the bright side, though, the system will be available to all users in that time [4].

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III. FAULT TOLERANCE IN ORACLE DATABASES

In this chapter, we discuss a specific technology, the Oracle Data Guard, for database fault tolerance. Before discussing the specific features and uses of Oracle Data Guard, we first provide an explanation of the product's intended purpose and value.

A. WHY ORACLE DATA GUARD

According to the official definition from the Oracle white paper, "Oracle Data Guard is the management, monitoring, and automation software infrastructure that creates, maintains, and monitors one or more standby databases to protect enterprise data from failures, disasters, errors, and corruptions" [21].

The same Oracle white paper also states:

Data Guard maintains these standby databases as synchronized copies of the production database. These standby databases can be located at remote disaster recovery sites thousands of miles away from the production data center, or they may be located in the same city, same campus, or even in the same building. If the production database becomes unavailable because of a planned or an unplanned outage, Data Guard can switch any standby database to [take on] the production role, thus minimizing the downtime associated with the outage, and preventing any data loss [22].

In addition to Oracle Data Guard, Oracle has another option, namely Oracle Active Data Guard, which is a superset of Oracle Data Guard. Oracle Active Data Guard is an option for the Oracle Database Enterprise Edition. It provides advanced capabilities, such as Data Guard functionality, including Real-Time Query, Automatic Block Repair, Far Sync, RMAN Block Change Tracking, Active Data Guard Rolling Upgrade, Global Database Services, and Application Continuity [23].

The ultimate goal of using Oracle Data Guard capability is to ensure the High Availability, Quality of Service, Data Protection, and Disaster Recovery [24]. It can be clearly deduced that Oracle Data Guard is a means of ensuring Fault Tolerance in the databases. To understand how particular characteristics of Oracle Data Guard meet the Fault Tolerance Objectives, it is important to recall the significant requirements described in the previous chapter: dependability, availability, reliability, safety, and maintainability. Reliability is ensured in Oracle Data Guard by real-time data protection. According to [19], "Oracle Data Guard enables zero data loss disaster recovery (DR) across any distance without impacting database performance. It repairs physical corruption without impacting availability" and saves network bandwidth without special-purpose network devices [25].

Availability is ensured in Oracle Data Guard by providing continuous access to data. According to [25], "It is essential when very little or no downtime is acceptable to perform maintenance activities. Activities, such as moving a table to another location in the database or even adding CPUs to hardware, should be transparent to the user in a high availability architecture" [25]. Oracle documentation [25] continues that, "Availability is the degree to which an application, service, or function is accessible on demand. Availability is measured by the perception of an application's user. If a user cannot access the system, it is said to be unavailable. Generally, the term downtime is used to refer to periods when a system is unavailable" [25]. Oracle documentation [25] also argues that, "Users who want their systems to be always ready to serve them need high availability. A system that is highly available is designed to provide uninterrupted computing services during peak periods, during most hours of the day, and most days of the week throughout the year; this measurement is often shown as 24x365. Such systems may also need a high availability solution for planned maintenance operations such as upgrading a system's hardware or software" [25].

Dependability and maintainability are ensured in both Oracle Data Guard and Oracle itself by constant software updates and customer support.

B. ORACLE DATABASE 12C

Oracle Database 12c has a single software application that is capable of serving multiple customers. This is referred to as multitenant architecture. Multitenant architecture enables managing databases as a cloud service. The main focus of the Oracle Database 12c is on efficiency, performance, security, and availability. Oracle Database 12c has two editions, which are the Enterprise Edition and the Standard Edition 2 [26].

1. An Overview

Oracle Database 12c offers many new concepts to its users. According to the official Oracle white paper [27], "Oracle Database 12c supports a new architecture that lets you have many 'sub databases' inside a single 'super database' ... The 'super database' is the multitenant container database—abbreviated as CDB; and the 'sub database' is the pluggable database—abbreviated as PDB. In other words, the new architecture lets you have many PDBs inside a single CDB" [27]. The new architecture is referred to as the multitenant architecture.

2. Database

Referring to Figure 6 provides a consolidated view into the inner workings of Oracle Database 12c. This screenshot is taken from Fernandez's book [28]. The software tool that was used to generate the screenshot is Spotlight [28]. This figure gives an insight into many topics that will be discussed in following sections, such as System Global Area (SGA), Redo logs, Achieve Logs, and Database Files [28].

Spotlight™on Oracle® - ORCL12C _ | C | X | File Yew Tools Help 作 浄 生・サ・エウや み は 2011年年間・6 1回の 『中山山山 Server Processes 0.00 blks/s 2594 2h 1m PGA Target 100 MB Container database Used 73.2 MI Buffer Cache 76.0 MB 4,970 blks/ Recycle Pool 0.00 XII Background Keep Pool 0.00 KB **Total Users 7** Dedicated Buffer Cache HII Ratio Flash Cache 0.00 b&x/s 1,640 KB/s LIMI chang Shared 4.69 68 Current SGA 139 MIII Total CPU Usage 0.00 10/6 Java Pool 4.00 Mil Large Pool 12.0 MD Parallel Query Redo Log Writer 10 Rado Buller 4.04 MB TGWR1 Groups 1,841 Parallel Exec Queu Shared Pool 176 MB Job Queue Shared Pool Used 0.00 KBA Result Cache EM Agent Size 0.00 KT Free 16.3 MB Finds/s 0.00 Total 19 lute On (AMM)

Figure 6. The Spotlight Tool's Representation of Oracle Database 12c

Source: I. Fernandez, *Beginning Oracle Database 12c Administration: From Novice to Professional*, 2nd ed., New York, NY: Apress IOUG, 2015, p. 48.

a. Database vs. Instance

In Oracle terminology, the term of "database" means the physical storage of data and information, and the term of "instance" means the software that is running on the server. The instance provides access to the data and information in the database and the resources that the software uses [29]. Another way to distinguish the instance from the database is to consider that the instance runs on the computer or server, and the database resides on hard drives that are attached to the server or computer [29].

The relationship between an instance and a database is presented in Figure 7.

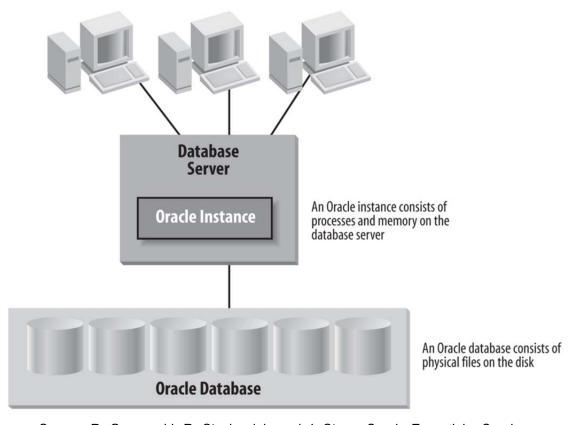


Figure 7. An Instance and a Database in the Same Figure

Source: R. Greenwald, R. Stackowiak, and J. Stern, *Oracle Essentials: Oracle Database 12c*, 5th ed. Sebastopol, CA: O'Reilly Media, Inc., 2013, p. 40.

The database is a physical structure. It consists of many files stored on physical disks. On the other hand, the instance is logical. The instance has processes on the server [29]. An instance is able to reach only one database. However, there can be more than one instance in the same database [29].

Oracle has two types of memory areas, the System Global Area (SGA) and Program Global Area (PGA). The System Global Area (SGA) is an area of shared memory. The Program Global Area (PGA) is an area of private memory for each process.

b. Data Files and Tablespaces

Data files are collections of files, which are grouped into tablespaces logically. They have descriptive names, such as DATA, INDEX, UNDO, and TEMP, in order to indicate the purposes of those files accordingly [28].

All data files have to be stored in tablespaces. A tablespace is a logical structure. Each tablespace has some physical structures known as data files. Each tablespace has to have one or more data files. Each data file has to belong to only one tablespace [29].

c. Control Files

Database control files have the locations of other physical files, which are data files and redo log files, that form the database. Database control files store very important information about the database such as [29]:

- The name of the database
- When the database was created
- Names and locations of data files and redo log files
- Tablespace information
- Data file offline ranges
- The log history and current log sequence information
- Archived log information
- Backup set, pieces, data file, and redo log information
- Data file copy information
- Checkpoint information [29]

The database control file, which contains the database startup settings, has two versions that are a pfile (text version) and an spfile (binary version). A pfile and an spfile have important information such as the memory that the system will use during its operation [28].

Listener.ora is another type of configuration file, which controls the operation of listener, a process in the database session [28].

Figure 8 shows the relationship among control files, data files, and redo log files.

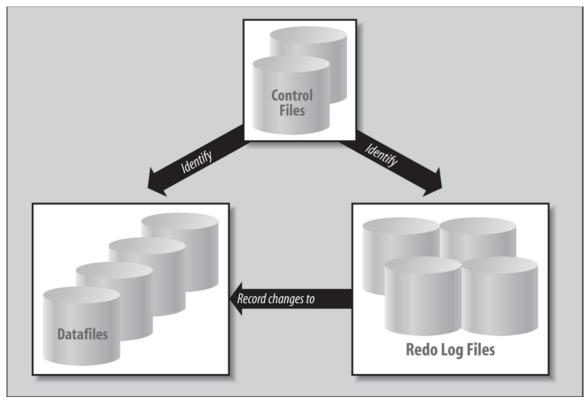


Figure 8. Data files, Redo Log files, and Control Files

Source: R. Greenwald, R. Stackowiak, and J. Stern, *Oracle Essentials: Oracle Database 12c*, 5th ed. Sebastopol, CA: O'Reilly Media, Inc., 2013, p. 43.

d. Redo Log Files

Redo log files hold a log of changes made to databases [29]. Those changes are results of transactions and internal Oracle activities.

From the fault tolerance perspective, redo log files are vitally important. When an instance failure occurs, some changes that were made to the database may not be applied. Redo log files are helpful at those moments; they can be used to play back the changes and apply them to the database. By doing that,

redo log files protect the database by having a fault tolerance to instance failures [29].

Moreover, redo log files are used for "undo" operations, which are the following operations after "Rollback." "Rollback" is basically going back to the last commit of the database [29].

C. ORACLE DATA GUARD

Oracle Data Guard, its configurations, advantages, and the inner workings of this technology are discussed in this section.

1. An Overview

According to the official definition from the Oracle white paper, "Oracle Data Guard is the management, monitoring, and automation software infrastructure that creates, maintains, and monitors one or more standby databases to protect enterprise data from failures, disasters, errors, and corruptions" [21].

The same Oracle white paper adds that

Data Guard maintains these standby databases as synchronized copies of the production database. These standby databases can be located at remote disaster recovery sites thousands of miles away from the production data center, or they may be located in the same city, same campus, or even in the same building. If the production database becomes unavailable because of a planned or an unplanned outage, Data Guard can switch any standby database to the production role, thus minimizing the downtime associated with the outage, and preventing any data loss [22].

2. Oracle Data Guard Configurations

According to [30], "An Oracle Data Guard configuration can contain one primary database and up to 30 destinations. The members of an Oracle Data Guard configuration are connected by Oracle Net." Moreover, Oracle Data Guard members can be remote from each other; they do not have to be in the same location [30].

Oracle Net and Oracle Net Services provide the connectivity across a distributed Oracle Data Guard configuration.

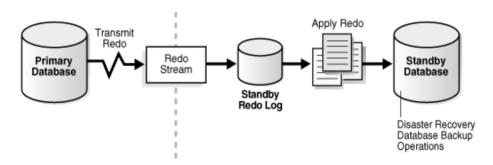
According to Oracle Database Online Documentation [31], "Oracle Net Services provides enterprise wide connectivity solutions in distributed, heterogeneous computing environments. Oracle Net Services ease the complexities of network configuration and management, maximize performance, and improve network diagnostic capabilities" [31]. The same documentation defines Oracle Net as "a component of Oracle Net Services, [which] enables a network session from a client application to an Oracle Database server" [31].

The most important point in terms of the members of an Oracle Data Guard configuration is they have to communicate each other, wherever they are. Thus, there is no limitation or no restriction on where they are physically located at [30].

The Oracle Data Guard configuration typically has one primary database, and up to 30 standby databases [30]. Standby databases can be any of these three: physical standby databases, which are a "physically identical copy of the primary database, with on disk database structures that are identical to the primary database on a block-for-block basis" [30]; logical standby databases, which have the "same logical information as the production database"; or snapshot standby databases, which are very much like physical or logical standby databases [30]. This third type of database "receives and archives redo data from a primary database" [30].

Figure 9 shows a typical Oracle Data Guard configuration. In the configuration shown, the primary database transmits redo data, which contains a recording of changes made to the database, to the standby database [30].

Figure 9. Typical Oracle Data Guard Configuration



Source: *Introduction to Oracle Data Guard.* (n.d.). Oracle. [Online]. Available: https://docs.oracle.com/database/121/sbydb/concepts.htm#sbydb4701.

Accessed: Feb.02, 2016.

3. Advantages of Oracle Data Guard

Oracle Data Guard has many advantages. According to [24], these advantages are:

- Disaster recovery, data protection, and high availability,
- Complete data protection,
- Efficient use of system resources,
- Flexibility in data protection to balance availability against performance requirements,
- Automatic gap detection and resolution,
- Centralized and simple management,
- Integration with Oracle Databases, and
- Automatic role transitions.

Out of all these advantages, the most noticeable of all are disaster recovery, data protection, and high availability. The data protection benefit also shines out in terms of reliability.

4. How Data Guard Synchronizes Standby Databases

Primary and standby databases in Oracle Data Guard configuration use Transmission Control Protocol/Internet Protocol (TCP/IP) to communicate using Oracle Net Services [23].

According to [23], "Data Guard automatically synchronizes the primary database and all standby databases by transmitting primary database redo logs" and applying that information to the standby databases [23].

a. Transport Services

Data guard transport services are responsible for transferring redo log files from the primary database to standby databases. This can be done automatically by the system itself, or it can be initiated manually by users. The changes that are made to the primary database are written to a local online log file via redo log files. Those redo log files are generated as users commit transactions, which are made at a primary database. Transport services rapidly transfer those redo log files from the primary database log buffer to standby databases [30]. There are two main reasons why this process is very efficient:

Firstly, Data Guard makes a direct transfer from memory, where the redo log files are allocated and sitting within the system global area, to standby databases. Therefore, this approach prevents disk input/output overhead on a primary database [30].

Secondly, Data Guard transmits database redo log files only. Compared to the other available technology for data protection over physically separate environments, which is storage remote-mirroring, Data Guard offers a much more light-weight approach. Recent Oracle tests have shown that Oracle Data Guard transfers up to 7 times less network volume and 27 times fewer network input/output operations [30]. Data Guard physical standby prevents input/output overhead by having logical replication instead of physical replication [30].

Remote-mirroring occurs in two ways, synchronous and asynchronous [30]. Synchronous remote-mirroring is accomplished by staying in communication with the remote system and waiting for it to finish its writing process. The recovery point objective of synchronous remote-mirroring is zero data loss, and the recovery time objective is seconds to minutes. Asynchronous remote-mirroring, on the other hand, is a store and forward technique [30]. It is mostly used for cases where there are longer distances between the primary system and remote systems. On the bright side, both synchronous and asynchronous remote-mirroring have minimal to zero data loss risk, and they both have a quick data recovery advantage [32]. However, they are expensive, and they require a lot of disk space, because the all of the changes that are made to the primary system must be transmitted as blocks [30], [32].

Data Guard has two transport services: as synchronous and asynchronous transport services, and synchronous Data Guard transport services have two options in Oracle Database 12c: Fast Sync and Far Sync [30].

Fast sync enables standby databases to inform the primary database as soon as they received redo logs. This option improves the overall performance by decreasing the time for the primary system, waiting for standby system to finish its commit process [30].

Far sync facilitates zero data loss to standby databases when there is failover [30]. This is done by making primary database to have the acknowledgement after successfully commit process in the standby databases. This option is available in Oracle Active Data Guard [30].

Asynchronous transport services are very much like store and forward method. In asynchronous transport services, the primary database simply does not wait for standby databases' acknowledgements for commit [30].

b. Redo Apply Services

According to [30], "Redo Apply services run on a physical standby database. Redo Apply reads redo records from a standby redo log file, performs Oracle validation to ensure that redo is not corrupt" [30]. After that, Redo Apply services apply redo changes to the standby database. Redo apply functions are independent from transport services [30].

c. Continuous Oracle Verification

Data Guard validates all redo logs before they are committed to the standby database. Data Guard also detects corruptions that occur because of lost-writes. Lost-writes primarily result from miscommunication between the database system and the hard drive. These corruptions may spread from standby databases to the primary one, or vice versa. Data Guard prevents further corruptions by performing lost-write validation at standby databases [30].

5. Managing the Data Guard Configuration

The most convenient way to interact with Oracle Database 12c is using the SQL*Plus application. The SQL*Plus application is available for various operating systems, including Windows, MAC OS X, and Linux.

Oracle Data Guard can be managed by Data Guard broker, which "automates and centralizes the creation, maintenance, and monitoring of a Data Guard configuration" [23]. Database administrators (DBAs) can use either Data Guard broker's command line interface or Oracle Enterprise Manager Cloud Control for interacting the Data Guard broker [23].

Oracle Enterprise Manager Cloud Control is a key product that the Oracle company uses primarily to provide a complete and integrated enterprise cloud management solution for information technology management [33]. DBAs use Oracle Enterprise Manager Cloud Control product via control console. A screenshot of the Oracle Enterprise Manager Cloud Control console appears in Figure 10.

Setup ▼ Help ▼ 🙎 GUEST_SUPER_ADMINI Log Out 🔘 ORACLE Enterprise Manager Cloud Control 12c 🌍 Grid 🕶 👩 Iargets 🕶 🏫 Eavorites 🕶 🙆 History 🕶 Search Target Name ▼ **Enterprise Summary** Page Refreshed Jul 29, 2011 4:39:05 AM PDT C ✓ Overview View All Targets ✓ Inventory and Usage Targets Monitored 2905 Show Hosts See Details View ▼ Platform Targets with Status 2241 Platform Hosts OS Patches Enterprise Linux Server release 5.6 (Carthage) ■ Down(93) Enterprise Linux AS release 4 (October Update 8) 15 No Metric Collection Error(267) Enterprise Linux Server release 5.4 (Carthage) SunOS 5 No ■Agent Unreachable(68) Red Hat Enterprise Linux Server release 6.0 (Santiago) ٧ 3 No Status Pending(338) Up(1,475) ✓ Compliance Summary Compliance Frameworks Compliance Standards Incidents View ▼ View Trends Open 1097 Updated in last 24 hours 696 Average Compliance Score (%) Category 🖨 🔞 🛕 🏲 O A P O A 4 Availability 95 244 60 No data to display Performance Security 122 Others 934 View + Problems Open 36 Without Service Request 36 Target Type Target Name Evaluations Updated in last 24 hours 30 **③** ▲ 🏲 **③** ▲ 🎄 example1.com 51 Suspended Executions (last 7 days) 15 example2.com Host 51 Problem Executions (last 7 days) 4267 X example3.com 51 Host 1 0 0 1 0 0 example4.com Action Required Executions (last 7 days) 0 ✓ Host example5.com 100 🗸 < → Patch Recommendations View by (Classification (Target Type ✓ Service Requests Other Recommendations My Oracle Support 0 10 20 30 40 50 60 You cannot access My Oracle Sunnort while in offline mode

Figure 10. A Picture of Oracle Enterprise Manager Cloud Control Console

Source: Overview of Oracle Enterprise Manager Cloud Control 12c. (n.d.). Oracle. [Online]. Available: http://docs.oracle.com/cd/E24628_01/doc.121/e25353/overview.htm#EMCON110. Accessed: Feb.09, 2016.

Oracle Enterprise Manager Cloud Control has setup assistants to enable creating Data Guard configurations. These assistants also simplify the process for DBAs [23]. The Data Guard management page, which is accessible on Oracle Enterprise Manager Cloud Control, can be seen in Figure 11.

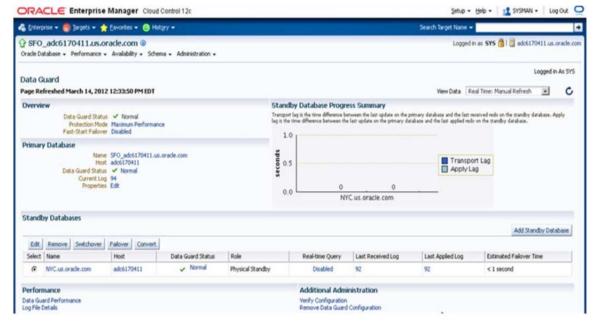


Figure 11. Data Guard Management Page

Source: Oracle Active Data Guard Real-Time Data Protection and Availability. (2015, Jan). Oracle. [Online]. Available: http://www.oracle.com/technetwork/database/availability/active-data-guard-wp-12c-1896127.pdf.

a. Switchover and Failover

There are two Data Guard role management services, such as switchover and failover [23]. The primary difference between a switchover and a failover is that the former is a planned event, whereas the latter is unplanned. A switchover is a planned event that is mainly used when there is a planned maintenance or software/hardware upgrade happening. The primary goal of switchover is to reduce the downtime, as well as to provide the zero data loss [23].

A failover assigns one of the standby databases as primary. In case of unplanned shutdown of the primary database, it does not require the standby database to be restarted before assuming the primary database's role [23]. There are two ways to initiate a failover. The first one is manual failover, where the DBA uses the Oracle Enterprise Manager Cloud Control Graphical User Interface (GUI), the Data Guard broker's command line interface, or SQL*Plus. The

second way is using fast-start Failover, where Data Guard provides an automatic failover [23].

b. Fast-Start Failover

Fast-start Failover enables [23] "Data Guard to automatically recover from a previously chosen standby database without requiring manual intervention to start the failover process" [23]. In fast-start Failover role, Data Guard also "continuously monitors the status of the configuration and initiates a failover if needed" [23].

c. Automating Client Failover

Although manual failover and fast-start failover are very important capabilities, they are not enough by themselves. In order to have an efficient high availability, all running processes and applications should be able to kill their existing transactional processes to or from the failed primary database, and direct themselves by quickly reconnecting to the new primary database [23].

According to [23], Data Guard broker can automatically assign a "standby database to the primary role," start "required database services appropriately for the primary role," and notify "application clients to disconnect from the failed primary database." Broker can also direct application clients to the new primary database, without requiring any further manual actions [23].

6. Tying Data Guard to Fault Tolerance

As discussed in Chapter II, there are four fault tolerance stages, which are error detection, damage confinement, error recovery, and fault treatment and continued system service.

Error detection is done by DGMGRL and Observer in the Data Guard Configuration. DGMGRL monitors the status of the configuration and tells whether the configuration is successful, DGMGRL gives details if the configuration is not successful. Observer monitors the fast-start failover

configuration and detects any errors in the primary database to perform an automatic failover to the standby database. Observer also automatically reinstates the primary database when it is up again.

Damage confinement is done and mitigated primarily by having redo log files, transport services, and flashback database in the Data Guard. Switchover and failover operations also contribute to the damage confinement stage.

Error recovery is provided by restoring or reinstating the failed database to an earlier state. The earlier state used for reinstatement is stored in the flashback database. The error recovery strategy in Oracle Data Guard is the backward recovery.

Fault treatment and continued system service is achieved by the combined efforts of redo log files, flashback databases, switchover and failover operations, and transport services. They all contribute to fault treatment and continued system service stage. However, the error removal from the system or databases requires manual interventions.

D. ORACLE VM VIRTUALBOX

Oracle VM Virtualbox and its capabilities are discussed in this section.

1. An Overview

"Oracle VM VirtualBox is a cross-platform virtualization application" [34]. It is installed on the existing operating system. It also boosts the capabilities of the existing computer "so that it can run multiple operating systems at the same time" [34]. There is no limitation on the number of virtual machines to run using Oracle VM VirtualBox, except the disk space and the memory [34].

2. Capabilities and Technical Aspects

Here are the main features of Oracle VM VirtualBox application:

 Portability: According to [34] "VirtualBox runs on a large number of 32-bit and 64-bit host operating systems, such as Windows, Mac, Linux, and Solaris." Oracle VM VirtualBox is a "hosted" hypervisor

- (type 2). The overview in [34] continues that Oracle VM VirtualBox "requires an existing operating system to be installed. It can thus run alongside existing applications on that host" [34]. Oracle VM VirtualBox is "functionally identical on all of the host platforms, and the same file and image formats are used. This allows you to run virtual machines created on one host on another host with a different host operating system; for example, a virtual machine that is created on Mac, can also run under Linux" [34].
- Compatibility with the existing hardware capabilities of the host computer: Users can use Oracle VM VirtualBox with some old hardware, where new features are not present [34].
- Extended hardware support: Oracle VM VirtualBox supports guest multiprocessing, USB device support, hardware compatibility with the existing hardware capabilities of the host computer, full Advanced Configuration and Power Interface (ACPI) support, multiscreen resolutions, built-in Internet Small Computer System Interface (iSCSI) support, and Preboot Execution Environment (PXE) Network boot [34].
- Remote machine display: The Oracle VM VirtualBox Remote Desktop Extension (VRDE) offers "high-performance remote access to any running virtual machine. This extension supports the Remote Desktop Protocol (RDP) originally built into Microsoft Windows, with special additions for full client USB support" [34].

Oracle VM VirtualBox runs on the following operating systems [34]:

- Windows hosts:
- Windows Vista SP1 and later (32-bit and 64-bit)
- Windows Server 2008 (64-bit)
- Windows Server 2008 R2 (64-bit)
- Windows 7 (32-bit and 64-bit)
- Windows 8 (32-bit and 64-bit)
- Windows 8.1 (32-bit and 64-bit)
- Windows 10 RTM build 10240 (32-bit and 64-bit)
- Windows Server 2012 (64-bit)

- Windows Server 2012 R2 (64-bit)
- Mac OS X hosts (64-bit):
 - 10.8 (Mountain Lion)
 - 10.9 (Mavericks)
 - 10.10 (Yosemite)
 - 10.11 (El Capitan)
- Linux hosts (32-bit and 64-bit):
 - Ubuntu 10.04 to 15.04
 - Debian GNU/Linux 6.0 ("Squeeze") and 8.0 ("Jessie")
 - Oracle Enterprise Linux 5, Oracle Linux 6 and 7
 - Redhat Enterprise Linux 5, 6 and 7
 - Fedora Core / Fedora 6 to 22
 - Gentoo Linux
 - openSUSE 11.4, 12.1, 12.2, 13.
 - Mandriva 2011
- Solaris hosts (64-bit only):
 - Solaris 11
 - Solaris 10 (U10 and higher)

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IV. DESIGNING AND IMPLEMENTING A FAULT TOLERANT DATABASE SYSTEM

A. DESIGN AND RATIONALE

The system that we designed has:

- Two databases, namely ORA12C and ORA12CX;
- A database administrator, who is using Data Guard command line interface (DGMGRL);
- An observer, which monitors two databases and performs a faststart failover when it is needed.

The general picture of the system is given in Figure 12.

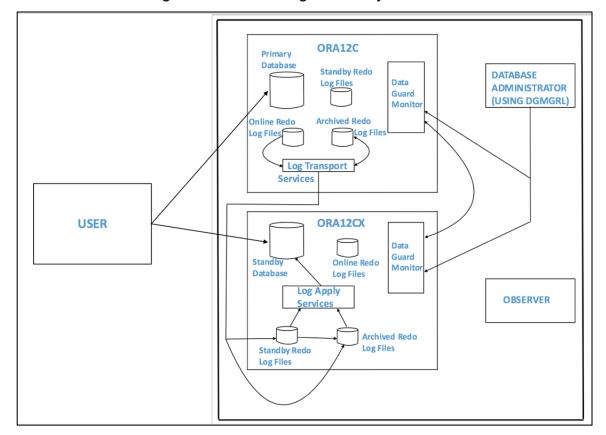


Figure 12. The Design of the System

The assumptions for the system that we designed derive from the actual needs of C2 systems in naval warfare ships. According to the *U.S. Department of Defense Dictionary of Military and Associated Terms*, the C2 system provides "...facilities, equipment, communications, procedures, and personnel essential for a commander to plan, direct, and control operations of assigned and attached forces pursuant to the missions assigned" [35].

C2 systems have to meet two minimum requirements, which are establishing and keeping the secure communication with the remote command authorities, and providing effective and correct supportive information to the captain and the officers in charge in the combat information center on a ship in a timely manner.

The design of the research system is focused on meeting the needs of the ship's interior information system that consistently supports the C2 system. That is the main reason why we chose to store the databases and other elements in the same operating system.

Our system uses the primary database, which is ORA12C, when everything is normal and there is no failure in the system. This means, ORA12C is the primary means of storing the information that supports the C2 system. When there is no fault (i.e., the primary database is up, running, and accessible to the other services), data will directly go to the primary database and our system will update the standby database in real time. Since both the primary and the standby databases are synchronized with the help of Oracle Data Guard architecture, the system itself will offer the redundancy by having the same data in multiple databases and places in the disk.

Having a problem in the standby database will not affect the actual transaction on the primary database and in the C2 system. The database managers will be informed of the problem in the standby database, and they can work on the problem without interfering with the active primary database. Oracle

Data Guard will sense and start storing changes in the primary database in redo log files until the future time when the standby database will be up and running.

When there is a problem encountered in the primary database, the system will perform a switch to the standby database. The standby database will take the responsibility of the primary database from that moment on until the database administrator fixes the problem in the primary database. After the standby database alters its role as primary, all changes will be stored on the standby database.

The database administrator can perform a switchover between two databases when there is a need. The need for switchover can be:

- Planned maintenance in the primary database;
- Planned maintenance on the primary database's hardware;
- Renewal of network cables in the primary database;
- Training purposes.

The database administrator can also perform a failover between two databases. The motives for failover can be:

- An unexpected problem or failure occurs in the primary database;
- An unexpected problem occurs on the primary database's hardware;
- A connection is lost or a network failure occurs;
- A need for training requires it.

The Data Guard is also capable of performing a failover when there is no available database administrator on duty, or to limit the down time of the system before the failover that is performed by the database administrator. The observer's primary job is to monitor the activities of the system and intervene when needed.

B. IMPLEMENTATION

We divided the implementation process of the system into nine main steps. Figure 13 shows the flow chart of the system implementation steps.

Installing Fedora Performing Switchover from the operating system on primary database to the standby VirtualBox database Installing Oracle Database Enabling flashback database 12c on Fedora Preparing the primary Performing manual failover from database and Creating the the primary database to the Physical Standby Database standby database **Enabling and using Fast-Start** Configuring Data Guard Broker Failover Changing Transport Mode using Broker properties and Changing **Protection Mode to Maximum** Availability

Figure 13. The Flow Chart of the System Implementation

1. Implementation Steps

First, we selected the Fedora operating system as the working environment because it is a stable and secure version of Linux. Furthermore, the Fedora operating system is a readily available operating system, which is free and open-source licensed. We also picked the VirtualBox application because it is a free hypervisor capable of hosting many operating systems, including Fedora. After setting up the initial environment, we started the second step of the implementation, which was installing Oracle Database 12c on Fedora operating system. The details of installing Oracle Database 12c on Fedora operating system is explained in the Appendix.

The third step was preparing the primary database and creating the physical standby database. According to [30], "physical standby database is kept synchronized with the primary database by using Redo Log files." We modified the third that the primary database in order to enable the primary database to talk to the standby database. After creating required directories, the password file, and the initialization parameter file for the standby database, we started up the standby instance and then ran the RMAN script for duplicating the primary to the standby.

The fourth step was configuring the Data Guard Broker. Data Guard Broker is a "framework that automates and centralizes the creation, maintenance, and monitoring the Data Guard configurations" [36]. This manages the entire Data Guard configuration including redo transport services, log apply services, switchovers, and failovers. In this step, we added DGMGRL listeners for each of the databases by using netmgr. We added ORA12C as the primary database and ORA12CX as the physical standby database in the DGMGRL configuration.

The fifth step was changing transport and protection modes. We changed the transport mode to synchronous using the Broker properties. Data Guard Transport Services are responsible for transferring redo log files from the primary to the standby database. We selected synchronous redo transport mode in our

system to make sure that the primary database waits for confirmation from the standby database when it sends redo log files. We also picked the protection mode as maximum availability to ensure that the primary database waits until it receives acknowledgement for applying redo log files from the standby or the timeout threshold expires. By selecting this mode, we emphasized availability as the first priority and the second priority as zero data loss protection.

The sixth step was performing a switchover from the primary database to the standby database. The key advantage of the switchover operation is no matter which transport service and protection mode are used, it always offers a zero data loss transition from the primary database to the standby database.

The seventh step was enabling the flashback database. The flashback database allows the database administrator and Data Guard configuration to "reinstate a failed primary database as a standby database after a failover occurred" [36].

The eighth step was performing a manual failover from the primary database to the standby database. This step was done to show how our system would react to an unexpected failure in the primary database. The expected actions for the standby database are taking the responsibility of the primary role and keeping the system running. This step was done manually from the DGMGRL console.

The last step was enabling and using fast-start failover. As described by [36], this allows the Data Guard configuration to automatically perform a failover to a standby database without needing a manual input or action from the database administrator. This goal is achieved by using the Observer, which according to [36] "continuously monitors the status of fast-start failover configuration." The Observer also automatically reinstates the failed primary database after the primary database is up again [36].

The details for installing Oracle Database 12c on the Fedora operating system and for setting up and managing the Oracle Data Guard configurations are presented in the Appendix section.

C. TESTING THE PERFORMANCE

Testing configurations and tests performed are presented in this section.

1. Testing Configuration

The configuration for testing the performance of the system is given in Figure 14.

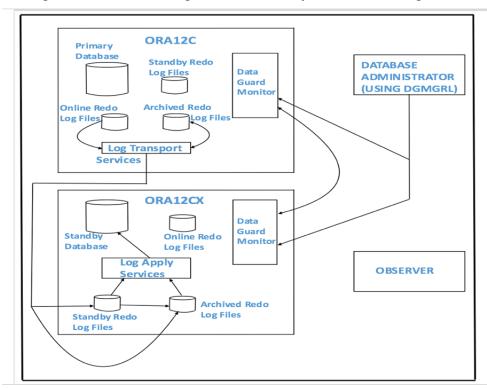


Figure 14. The Configuration of the System for Testing

2. Tests Performed

In order to test the capabilities of the system, numerous tests were performed and are described in the following subsections.

a. Connectivity

The connectivity test was performed to verify that both databases were up and running. In this test, we started up the listener, primary database (ORA12C), standby database (ORA12CX), and the observer before running Java codes. We also connected the primary database (ORA12C) using the DGMGRL to check the current configuration of the system.

We followed the following sequence of events during the connectivity testing scenario:

- 1. The primary database (ORA12C), the standby database (ORA12CX), and the observer were started.
- 2. The configuration of the system was checked by using DGMGRL.
- 3. The system configuration status was checked to confirm it as "SUCCESS."
- 4. Another terminal was opened prior to running the Java code for the connectivity testing.
- TestDBConnection.java code was run to check the connectivity of the primary database (ORA12C).
- 6. TestDBConnection2.java code was run to check the connectivity of the standby database (ORA12CX).

b. User Creation and Granting Roles

The user creation and granting roles test was performed in order to check what happens when the database administrator creates a user in the primary database, grants some roles to that user, and creates a table named test, by connecting to the system as that user. In this test, we started up the listener, primary database (ORA12C), standby database (ORA12CX), and the observer before running the test. We also connected the primary database (ORA12C) using the DGMGRL to check the current configuration of the system.

We followed the following sequence of events during the user creation and granting roles testing scenario:

- 1. The primary database (ORA12C), the standby database (ORA12CX), and the observer were started.
- 2. The configuration of the system was checked by using DGMGRL.
- 3. The system configuration status was checked to confirm it as "SUCCESS."
- 4. A new user, named "c##kadir," was created in the primary database in the SQL*Plus command line. Required grants were given to the user "c##kadir" in order to create a table, add to, update, and delete from the table.
- 5. A connection to the primary database (ORA12C) was established with the newly created user "c##kadir."
- 6. A new table, named "test," was created. One column was created for testing purposes in the "test" table. Five entries were inserted into the "test" table.
- 7. After completing the creating the user, granting roles, creating a table using that user's credentials, and inserting some values into that table, a new terminal was opened to connect to the standby database (ORA12CX).
- 8. A connection to the primary database (ORA12C) was established with the recently created user "c##kadir."
- 9. The "test" table and all entries were verified in the standby database (ORA12CX) using the user "c##kadir" credentials.

c. Reading from Databases

The reading from databases test was performed in order to confirm that the table count numbers in both databases were the same. In this test, we started up the listener, primary database (ORA12C), standby database (ORA12CX), and the observer before running the test. We also connected the primary database (ORA12C) using the DGMGRL to check the current configuration of the system.

We followed the following sequence of events during the reading from databases testing scenario:

1. The primary database (ORA12C), the standby database (ORA12CX), and the observer were started.

- 2. The configuration of the system was checked by using DGMGRL.
- 3. The system configuration status was checked to confirm it as "SUCCESS."
- 4. Another terminal was opened prior to running the Java code for the connectivity testing. The connectivity test was repeated.
- 5. Many entries were inserted in the "test" table.
- 6. TableCount.java code was run to check the number of entries of two databases.
- 7. The verification of the test was made by checking they both had the same number of entries. The result that has equal numbers in both tables and the elapsed times for reading from two databases can be observed in Figure 15.

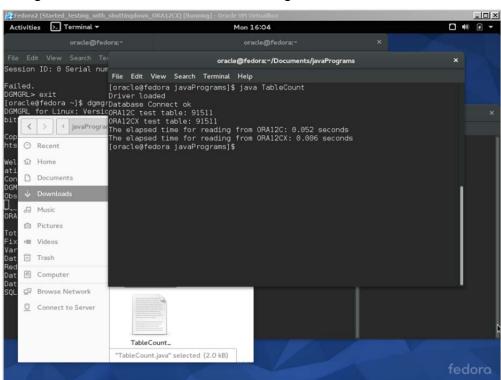


Figure 15. The Result of the Reading from Databases Test

d. Writing to Databases

The writing to databases test was performed in order to determine whether both databases synchronize with each other when a client makes

changes in the test table. In this test, we started up the listener, primary database (ORA12C), standby database (ORA12CX), and the observer before running the test. We also connected the primary database (ORA12C) using the DGMGRL to check the current configuration of the system.

We followed the following sequence of events during the writing to databases testing scenario:

- 1. The primary database (ORA12C), the standby database (ORA12CX), and the observer were started.
- 2. The configuration of the system was checked by using DGMGRL.
- 3. The system configuration status was checked to confirm it as "SUCCESS."
- 4. Another terminal was opened prior to running the Java code for the connectivity testing. The connectivity test was repeated.
- 5. UpdatingTable.java code was run to insert 2000 rows in the test table. In this step, there were no failures in the databases. The output of this step can be seen in Figure 16.
- 6. UpdatingTable.java code was run to insert 2000 rows in the test table again. In this step, there was an error injected to the standby database. The error was created by shutting down the standby database. The output of this step can be seen in Figure 17.
- 7. UpdatingTable.java code was run one more time to insert 2000 rows in the test table. The standby database was still inaccessible in this step. This step was done to insert more entries in the table to cause more transactional complexity. The output of this step can be seen in Figure 18.
- 8. The standby database (ORA12CX) was started again. Immediately after that, TableCount.java code was run to check the number of entries in both tables. The verification of the test is made by checking to see that they both have the same number of entries. The output of this step can be seen in Figure 19.

Figure 16. Inserting 2000 Rows When Both Databases Are Up

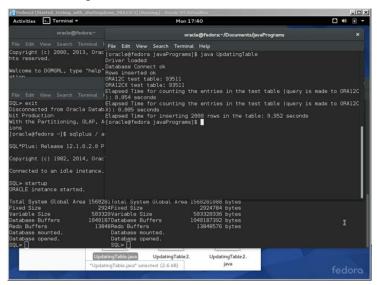


Figure 17. Inserting 2000 Rows When the Standby Database Is Down

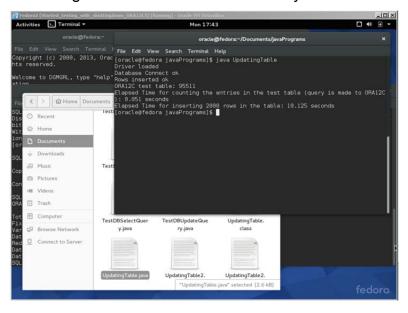


Figure 18. Inserting 2000 More Rows When the Standby Database Is Down

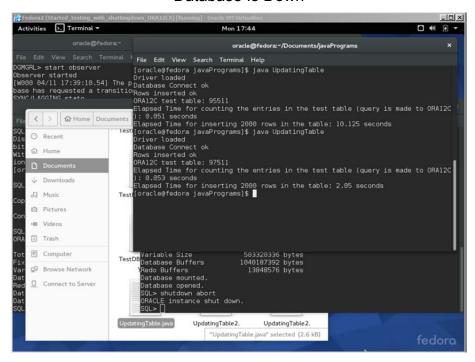
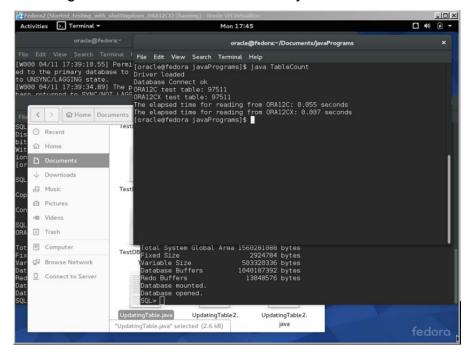


Figure 19. Two Databases Are Synchronized



e. Multiple Clients Trying to Access the System

The multiple clients trying to access the system test was performed in order to check what happens when two clients try to connect to the system and make changes in the test table. In this test, we started up the listener, primary database (ORA12C), standby database (ORA12CX), and the observer before running Java codes. We also connected the primary database (ORA12C) using the DGMGRL to check the current configuration of the system.

We followed the following sequence of events during the multiple clients trying to access the system testing scenario:

- 1. The primary database (ORA12C), the standby database (ORA12CX), and the observer were started.
- 2. The configuration of the system was checked by using DGMGRL.
- 3. The system configuration status was checked to confirm it as "SUCCESS."
- 4. Another terminal was opened prior to running the Java code for the connectivity testing. The connectivity test was repeated.
- 5. Before running the test, TableCount.java code was run to check the number of entries in the test table. The output of this step can be seen in Figure 20.
- 6. UpdatingTable.java code was run on two different terminals at the same time to insert a total of 4000 rows in the test table. In this step, there were no failures in the databases. The output of this step can be seen in Figure 21.
- 7. TableCount.java code was run one more time to verify that two numbers were matched. The output of this step can be seen in Figure 22.
- 8. Steps 5 was repeated one more time. In this step, there was an error injected into the standby database. The error was created by shutting down the standby database. UpdatingTable.java code was run on two different terminals at the same time to insert a total of 4000 rows in the test table when there was an error in the system. The output of this step can be seen in Figure 23 and Figure 24.
- 9. The standby database (ORA12CX) was started again. Immediately after that, TableCount.java code was run to check the number of

entries in both tables. The verification of the test is made by checking to see that they both have the same number of entries. The output of this step can be seen in Figure 25.



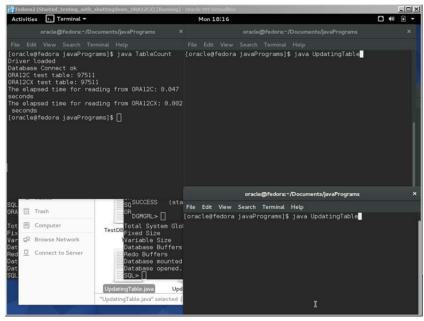


Figure 21. Inserting 4000 Rows from Two Different Terminals

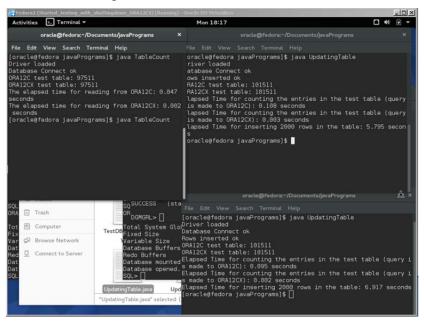


Figure 22. Two Databases Are Synchronized

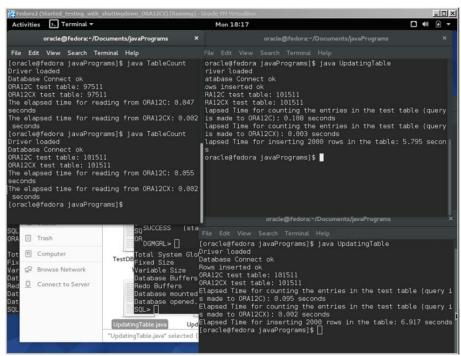


Figure 23. The Output Before Running Step 8

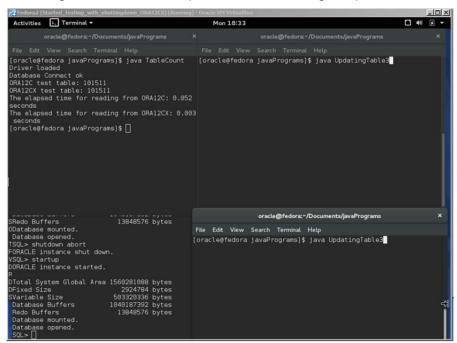


Figure 24. Inserting 4000 Rows from Two Different Terminals When the Standby Database Is Down

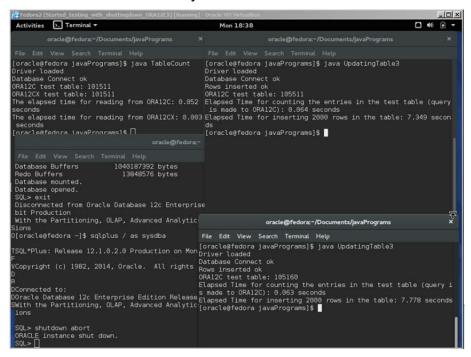
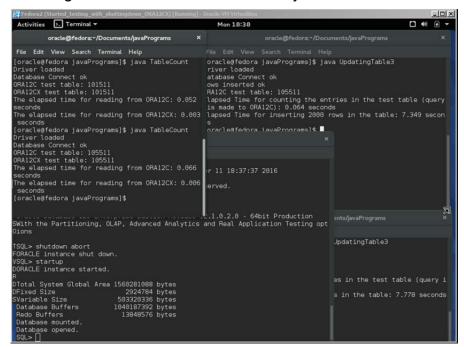


Figure 25. Two Databases Are Synchronized



f. Writing to Databases as Multiple Clients Are Trying to Access the System with Multiple Switchovers and Failovers

The writing to databases and multiple clients trying to access the system with multiple switchovers and failovers test was performed in order to check what happens when two clients try to connect to the system and make changes in the test table when there is more than one switchover and/or failover in the system. In this test, we started up the listener, primary database (ORA12C), standby database (ORA12CX), and the observer before running Java codes. We also connected the primary database (ORA12C) using the DGMGRL to check the current configuration of the system.

We followed the following sequence of steps during the writing to databases and multiple clients trying to access the system with multiple switchovers and failovers testing scenario:

- 1. The primary database (ORA12C), the standby database (ORA12CX), and the observer were started.
- 2. The configuration of the system was checked by using DGMGRL.
- 3. The system configuration status was checked to confirm it as "SUCCESS."
- 4. Another terminal was opened prior to running the Java code for the connectivity testing. The connectivity test was repeated.
- 5. Before running the test, TableCount.java code was run to check the number of entries in the test table. The output of this step can be seen in Figure 26.
- 6. UpdatingTable_Switchover_v3.java code was started running. The number of entries that were to be inserted into the test table was 10000.
- 7. While UpdatingTable_Switchover_v3.java code was running, we manually initiated a switchover to the standby database (switchover #1). In this step, ORA12CX became the primary database. The output of this step can be seen in Figure 27.
- 8. After the switchover was successfully performed, UpdatingTable_Switchover_v3.java code started running again. We again manually initiated a switchover to the standby database

- (switchover #2). In this step, ORA12C became the primary database. The output of this step can be seen in Figure 28.
- 9. We repeated the step 7 (switchover #3). The output of this step can be seen in Figure 29.
- 10. We repeated the step 8 (switchover #4). The output of this step can be seen in Figure 30.
- 11. UpdatingTable_Switchover_v3.java code finished running. The verification of the test is made by checking the number of entries in the system before and after running UpdatingTable_Switchover_v3.java code. We saw the number of entries was 10000 more in comparison to the previous number, which we had in Step 5. The output of this step can be seen in Figure 31.
- 12. The same steps 5 through 11 were run one more time with UpdatingTable_Switchover_v4.java code. In this version of the code, there were no printout statements. The verification of the test is made by checking the number of entries in the system before and after running UpdatingTable_Switchover_v4.java code. We saw the number of entries was 10000 more in comparison to the previous number before running the code. The output of this step can be seen in Figure 32.
- 13. After finishing the having multiple switchovers while inserting into the test table, we run the same test with two UpdatingTable_Switchover_v3.java codes running at the same time as multiple threads. Before running the test, TableCount.java code was run to check the number of entries in the test table. The output of this step can be seen in Figure 33.
- 14. Two UpdatingTable_Switchover_v3.java codes were started running at the same time. The number of entries to be inserted into the test table was 20000 in total.
- 15. While UpdatingTable_Switchover_v3.java codes were running, we manually initiated a switchover to the standby database (switchover #1). In this step, ORA12CX became the primary database. The output of this step can be seen in Figure 34.
- 16. After the switchover was successfully performed, UpdatingTable_Switchover_v3.java codes started running again. We again manually initiated a switchover to the standby database (switchover #2). In this step, ORA12C became the primary database. The output of this step can be seen in Figure 35.

- 17. We repeated the step 14 (switchover #3). The output of this step can be seen in Figure 36.
- 18. We repeated the step 15 (switchover #4). The output of this step can be seen in Figure 37.
- 19. UpdatingTable_Switchover_v3.java codes finished running. The verification of the test is made by checking the number of entries in the system before and after running UpdatingTable_Switchover_v3.java codes. We saw the number of entries was 20000 more compared to the previous number, which we had in Step 12. The output of this step can be seen in Figure 38.
- 20. The same steps 13 through 19 were run one more time with UpdatingTable_Switchover_v4.java code. In this version of the code, there were no printout statements. The verification of the test is made by checking the number of entries in the system before and after running UpdatingTable_Switchover_v4.java code. We saw the number of entries were 20000 more compared to the previous number before running the code. The output of this step can be seen in Figure 39.
- 21. After finishing the having multiple switchovers while inserting into the test table with multiple threads, we ran another test to evaluate the behavior of the system when there is a failover. Before running the test, TableCount.java code was run to check the number of entries in the test table. The output of this step can be seen in Figure 40.
- 22. Two UpdatingTable_Switchover_v3.java codes were started running at the same time. The number of entries to be inserted into the test table was 20000 in total. As they were running, we injected an error in the system by manually shutting down the primary database (ORA12C). The observer sensed the error in 30 seconds and initiated a failover to the standby database (ORA12CX). After the failover, the two threads started running once again. After waiting a couple of seconds, we manually started up the ORA12C instance. The observer started reinstating ORA12C as the two threads were running.
- 23. UpdatingTable_Switchover_v3.java codes finished running. The verification of the test is made by checking the number of entries in the system before and after running UpdatingTable_Switchover_v3.java codes. We saw the number of entries was 20001 more compared to the previous number, which we had in Step 19. The output of this step can be seen in Figure 41.

24. Lastly, we ran the same test with UpdatingTable_Switchover_v4.java code, which had no printout statements. The verification of the test is made by checking the number of entries in the system before and after running UpdatingTable_Switchover_v4.java codes. We saw the number of entries was 20001 more compared to the previous number before running the code. The output of this step can be seen in Figure 42.

Figure 26. The Number of Entries in the Test Table Before Running the Four Switchovers in a Single Thread Test

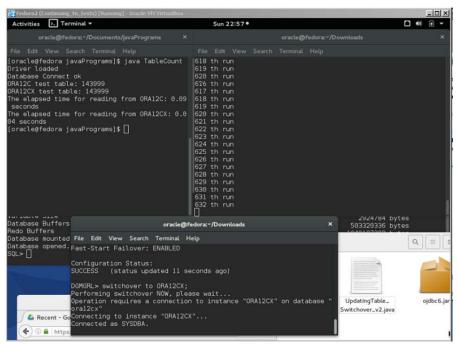


Figure 27. Switchover #1 (Single Thread Scenario)

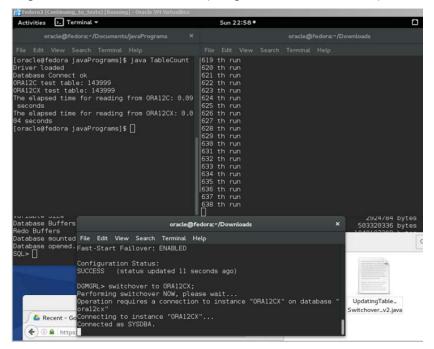


Figure 28. Switchover #2 (Single Thread Scenario)

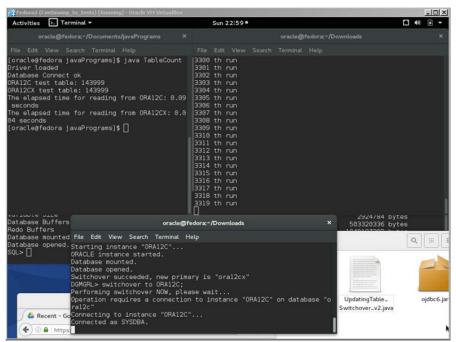


Figure 29. Switchover #3 (Single Thread Scenario)

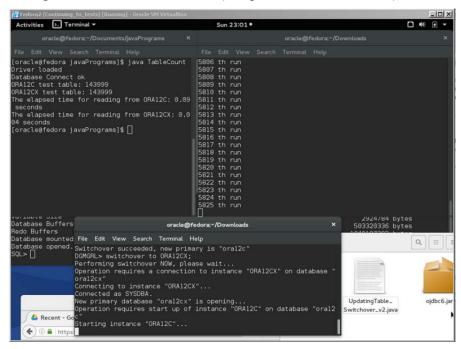
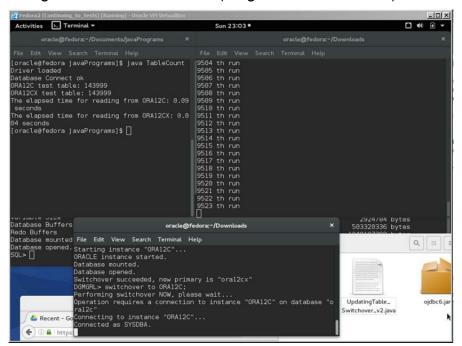


Figure 30. Switchover #4 (Single Thread Scenario)



Activities > Terminal -Sun 23:05 •

Figure 31. The End of Four Switchovers in a Single Thread Test

[oraclededora javaPrograms]\$ java TableCount Driver loaded Jatabase Connect ok DRAI2C test table: 143999 JRAI2CX test table: 143999 JRAI2CX test table: 143999 JRE elapsed time for reading from ORAI2C: 0.09 seconds ne elapsed time for reading from ORA12CX: 0.0 4 seconds pracle@fedora javaPrograms]\$ java TableCount seconds he elapsed time for reading from ORA12CX: 0.1 seconds oracle@fedora javaPrograms]\$ Q ||| Protection Mode: MaxAvailability Members: oral2c - Primary database oral2cx - (*) Physical standby database Fast-Start Failover: ENABLED UpdatingTable_ ojdbc6.ja Configuration Status: SUCCESS (status updated 103 seconds ago) witchover_v2.java & Recent - Go ♦ ① A https:DGMGRL>

Figure 32. The End of Four Switchovers in a Single Thread Test (No **Printout Statements**)

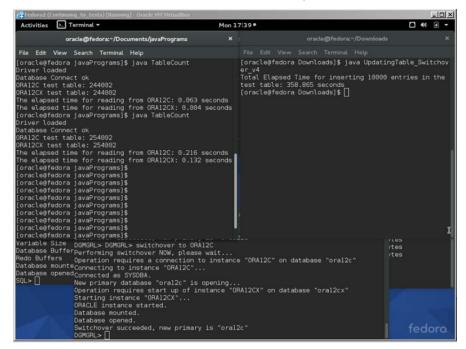


Figure 33. The Number of Entries in the Test Table Before Running the Four Switchovers in Multiple Threads Test

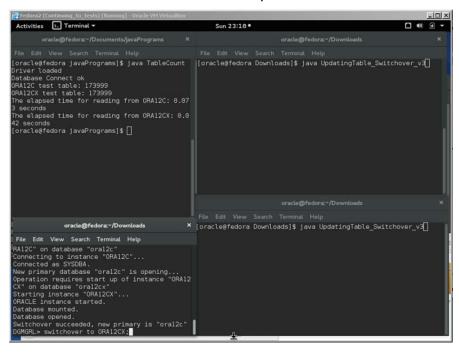


Figure 34. Switchover #1 (Multiple Threads Scenario)

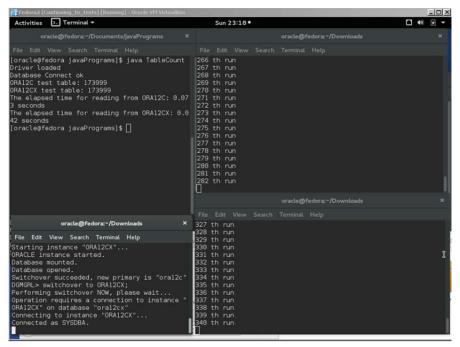


Figure 35. Switchover #2 (Multiple Threads Scenario)

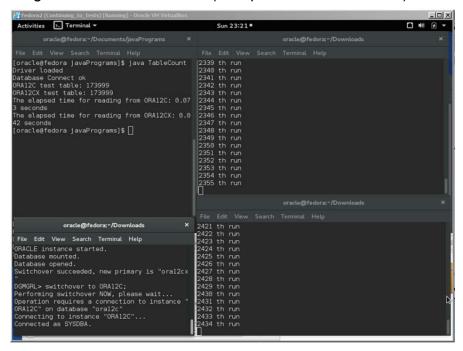


Figure 36. Switchover #3 (Multiple Threads Scenario)

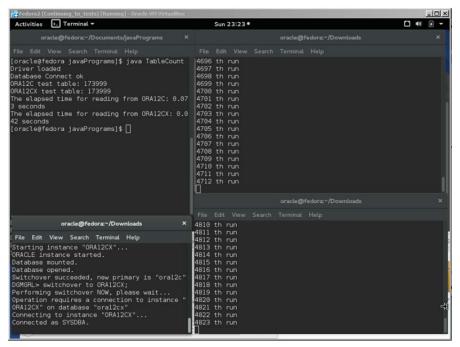


Figure 37. Switchover #4 (Multiple Threads Scenario)

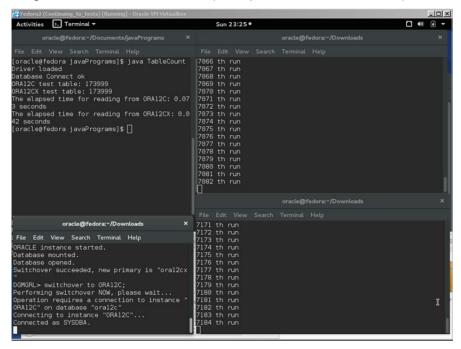


Figure 38. The End of Four Switchovers in Multiple Threads Test

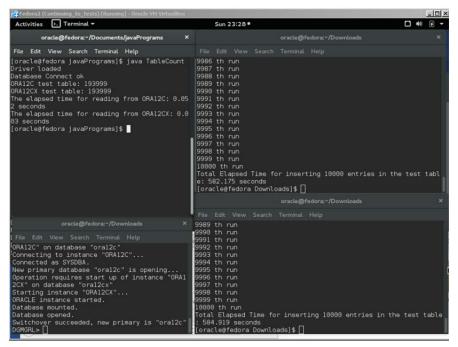


Figure 39. The End of Four Switchovers in Multiple Threads Test (No Printout Statements)

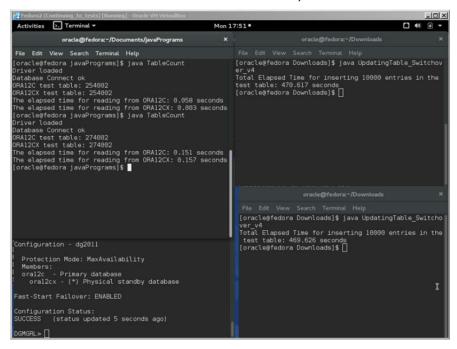
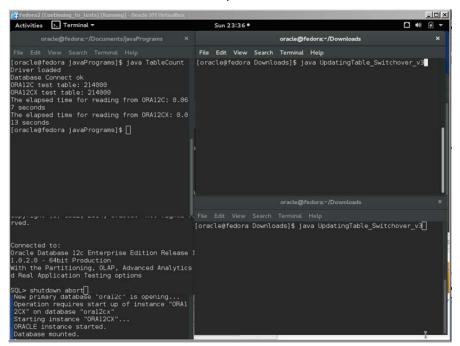


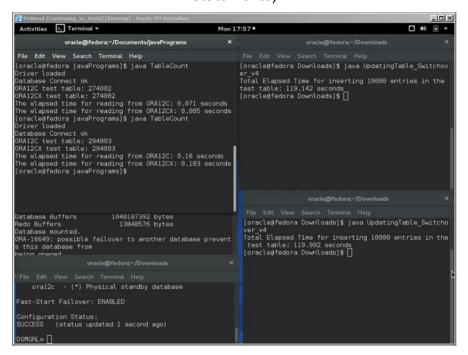
Figure 40. The Number of Entries in the Test Table Before Running a Failover in Multiple Threads Test



Activities Terminal T

Figure 41. The End of a Failover in Multiple Threads Test

Figure 42. The End of a Failover in Multiple Threads Test (No Printout Statements)



D. REVIEWING TEST RESULTS

There were six different tests performed in order to assess the functionality of the system. Those were connectivity, user creation and granting roles, reading from databases, writing to databases, multiple clients trying to access the system, and writing to databases as multiple clients are trying to access the system with multiple switchovers and failovers.

In the connectivity test, the expected outcome was that the system is up, running, and accessible. The expected result was met in this test.

In the user creation and granting roles test, the expected outcome was to see that when a user was created and granted some roles in one database, the same user was created in the other database. For this test, we created a user "c##kadir" in the primary database with permissions for creating a table, adding to it, updating, and deleting from the table. We also created a "test" table for that user. We observed that the same user was created in the other database with the "test" table. The expected result was met in this test.

In the reading from databases test, the expected outcome was to see the same number of entries in both databases' "test" tables. This test was performed when two databases were up, running, and accessible. We observed that after inserting some entries into the primary database's "test" table, the two databases synchronized with each other and the total number of entries in both "test" tables stayed the same. It took 0.052 seconds to read the number of entries in the "test" table from the primary database (ORA12C) and 0.006 seconds from the standby database (ORA12CX), respectively. The expected result was met in this test.

In the writing to databases test, the expected outcome was to see the same number of entries in both databases' "test" tables after inserting some entries in the primary database when the standby database was shut down. First, we inserted 2000 entries into the primary database's "test" table when the standby database was up. It took 9.9 seconds to insert 2000 entries. After inserting 2000 entries in the primary database's "test" table, we observed that

both databases had the same number of entries, which also meant that they were synchronized as expected. Second, we inserted 4000 entries into the primary database's "test" table when the standby database was shut down. It took 12 seconds to insert 4000 entries in the primary database's "test" table. After inserting 4000 entries in the primary database's "test" table, we started up the standby database. It took 11.5 seconds for the standby database to start up and get ready. After that, we observed that both databases had the same number of entries, which also meant that they were synchronized as expected. The expected result was met in this test.

In the multiple clients trying to access the system test, the expected outcome was to see the system responded to all clients that were trying to connect to databases and updating the "test" table. The expected outcome was also to see the same number of entries in both databases. In this test, we ran two Java codes that inserted 2000 entries each in the "test" table simultaneously from two terminals. We also performed the same test with a failure in the standby database. After inserting the total of 4000 entries in the primary database's "test" table, we observed that both databases had the same number of entries, which also meant that they were synchronized as expected. It took 8 seconds to insert 4000 entries in the primary database's "test" table. It took 11.5 seconds for the standby database to start up and get ready. The expected result was met in this test.

In the writing to databases as multiple clients are trying to access the system with multiple switchovers and failovers test, the expected outcome was to see the system responded to all clients that were trying to connect to databases and updating the "test" table. The expected outcome was also to see the same number of entries in both databases. In this test, we first ran one java code that inserted 10000 entries in the "test" table as there were four switchovers that happened in the system. It took 420 seconds to insert 10000 entries in the system as the four switchovers occurred. The same test was also run one more

time with no printout statements. This time, it took 358 seconds to insert 10000 entries in the system as four switchovers occurred.

Next, we ran two Java codes that inserted 20000 entries in the "test" table simultaneously from two terminals as there were four switchovers occurring in the system. It took 584 seconds to insert 20000 entries in the system as four switchovers occurred. The same test was also run one more time with no printout statements. This time, it took 470 seconds to insert 20000 entries in the system as four switchovers occurred.

Lastly, we ran two Java codes that inserted 20000 entries in the "test" table simultaneously from two terminals as one failover happened in the system. The observer was able to sense the error approximately 30 seconds after we manually shut down the primary database (ORA12C). It took 122 seconds to insert 20000 entries in the system as one failover occurred. The same test was also run one more time with no printout statements. This time, it took 119 seconds to insert 20000 entries in the system as one failover occurred. In the failover scenario, we observed that there was one more entry in the databases. This was a normal system behavior in maximum availability mode according to data protection definitions as described in *Data Guard Protection Modes* [30].

In all three scenarios in which multiple clients are trying to access the system with multiple switchovers and failovers, the expected number of entries was observed after the performed tests. The expected results were met in those tests.

We have made another observation regarding the switchover and failover durations. It took 45 seconds to perform a switchover and 75 seconds to perform a failover. This observation also agrees with *Client Failover Best Practices for Highly Available Oracle Databases* [37].

V. CONCLUSIONS AND FUTURE WORK

A. SUMMARY

The purpose of this thesis is to find an alternative way for ensuring the fault tolerance in C2 systems of naval platforms. It is highly critical for C2 systems of naval platforms to be robust to failures.

This thesis primarily focused on the databases of C2 systems, which provide all required data to operators, officers, and the captains in order to help them make decisions using their technical knowledge combined with detailed scientific information. The data that are stored in these databases can vary in a high spectrum. The data can simply be the velocity and the direction of the wind, which can be expressed by integers or floats, or be sophisticated information of the weapons on board.

In this thesis, we showed that there is a potentially cheaper way to provide a fault tolerant mechanism for C2 database systems. We designed and implemented a system using Oracle Database 12c Enterprise Edition and its features. Specifically, we used two databases in the system, ORA12C and ORA12CX. Then we connected those two databases with the Data Guard capability.

Six different tests were performed to check the capabilities and performance of the system. We observed that all expected results were met and the system behaved in accordance with the fault tolerance objectives.

Our experiments showed that there is a way to have a robust, fast, and fault tolerant system for storing databases of C2 systems by using a virtual environment rather than a physical redundancy. Our alternative approach also can decrease the cost of having multiple of redundant servers that store databases. This approach can be used in many fields such as small businesses and large enterprises that require a fault tolerant solution for physical, human-

made, design, and interaction faults. The military domain can also take advantage of the approach presented in this work.

B. FUTURE WORK

This thesis makes the initial effort to ensure fault tolerance by working in a virtual environment. The system that we designed here can be tested with the data flow in a real C2 system of a naval platform as a future work idea.

The framework in this thesis is stored in the same virtual machine by using the local machine's connections to simulate the C2 system database that is physically located on the ship. The next step can be installing and modifying a fault tolerant system in distant locations, thus enabling a redundant off-site database system that can function as a means of backup as well as a fault-tolerant node.

APPENDIX

A. INSTALLING ORACLE DATABASE 12C ON FEDORA OPERATING SYSTEM

In this appendix, we describe how we installed Oracle Database 12c on the Fedora operating system, which runs on the Virtual Box application. We primarily used the http://dbaora.com/install-oracle-12c-release-1-12-1-on-fedora-22/ website, which is included in the references section.

First, we downloaded the Oracle Database 12c software from the official Oracle website (http://www.oracle.com/technetwork/database/enterprise-edition/downloads/index.html). Before installing the software on the operating system, we made sure that the operating system was updated to the latest version properly by entering "dnf update" command in the command line.

The groups and users that were used for installing the database are given here:

"#groups for database management

groupadd -g 54321 oinstall

groupadd -g 54322 dba

groupadd -g 54323 oper

groupadd -g 54324 backupdba

groupadd -g 54325 dgdba

groupadd -g 54326 kmdba

groupadd -g 54327 asmdba

groupadd -g 54328 asmoper

groupadd -g 54329 asmadmin

#users for database management

useradd -u 54321 -g oinstall -G dba,oper,backupdba,dgdba,kmdba oracle"

After choosing the password for our database, we checked to see which packets were missing for the installation of the database by typing in the command line as follows:

"rpm -q --qf '%{NAME}-%{VERSION}-%{RELEASE}(%{ARCH})\n' binutils \compat-libstdc++-33 \gcc \gcc-c++ \glibc \glibc-common \glibc-devel \glibc-headers \ksh \libaio \libaio-devel \libgcc \libstdc++ \libstdc++-devel \libXext

\libXtst \libX11 \libXau \libXi \make \sysstat \unixODBC \unixODBC-devel \zlib-devel"

After figuring out the missing packets, we downloaded and installed them on the operating system. We added and applied the kernel parameters, which are recommended by Oracle, to the operating system. All of these given parameters were written into "/etc/sysctl.conf" file as follows:

"--kernel parameters for 12gR1 installation fs.file-max = 6815744 kernel.sem = 250 32000 100 128 kernel.shmmni = 4096 kernel.shmmall = 1073741824 kernel.shmmax = 4398046511104 net.core.rmem_default = 262144 net.core.rmem_max = 4194304 net.core.wmem_default = 262144 net.core.wmem_max = 1048576 fs.aio-max-nr = 1048576 net.ipv4.ip_local_port_range = 9000 65500 /sbin/sysctl -p"

The recommended shell limits for the Oracle user were written into "/etc/ security/limit.conf" file as follows:

"--shell limits for users oracle 12gR1 oracle soft nofile 1024 oracle hard nofile 65536 oracle soft nproc 2047 oracle hard nproc 16384 oracle soft stack 10240 oracle hard stack 32768"

In order to make the installation more user-friendly and more hassle-free, we disabled the SELINUX, which is "a reference implementation of the Flask security architecture for flexible mandatory access control. It was created to demonstrate the value of flexible mandatory access controls and how such controls could be added to an operating system" [38].

For the sake of ease of use, we created the bash profile for the "oracle" user to effectively use the shortcuts. The bash profile settings are provided here:

```
"# Oracle Settings
export TMP=/tmp
export ORACLE HOSTNAME=fedora.dbaora.com
export ORACLE UNQNAME=ORA12C
export ORACLE BASE=/ora01/app/oracle
export ORACLE HOME=$ORACLE BASE/product/12.1.0/db 1
export ORACLE SID=ORA12C
PATH=/usr/sbin:$PATH:$ORACLE HOME/bin
export LD LIBRARY PATH=$ORACLE HOME/lib:/lib:/usr/lib;
export CLASSPATH=$ORACLE HOME/ilib:$ORACLE HOME/rdbms/ilib:
alias cdob='cd $ORACLE BASE'
alias cdoh='cd $ORACLE HOME'
alias tns='cd $ORACLE HOME/network/admin'
alias envo='env | grep ORACLE'
umask 022
envo"
```

Following are the installation commands:

```
"##this will increase the permissions su
##unzipping the software
unzip linuxamd64_12102_database_1of2.zip
unzip linuxamd64_12102_database_2of2.zip
[oracle@fedora ~]$ alias envo cdob cdoh tns
alias envo='env | grep ORACLE'
alias cdob='cd $ORACLE_BASE'
alias cdoh='cd $ORACLE_HOME'
alias tns='cd $ORACLE_HOME/network/admin'
```

##runs alias command envo to display environment settings

```
[oracle@fedora ~]$ envo
ORACLE_UNQNAME=ORA12C
ORACLE_SID=ORA12C
ORACLE_BASE=/ora01/app/oracle
ORACLE_HOSTNAME=fedora.dbaora.com
ORACLE_HOME=/ora01/app/oracle/product/12.1.0/db 1
```

##runs alias command cdob and cdoh to check ORACLE_BASE, ORACLE HOME

[oracle@fedora ~]\$ cdob
[oracle@fedora oracle]\$ pwd
/ora01/app/oracle

[oracle@fedora db_1]\$ cdoh [oracle@fedora db_1]\$ pwd /ora01/app/oracle/product/12.1.0/db_1

##finally runs the installation

./runInstaller"

After this step, an installation window of the Oracle Database 12c was displayed on the screen. There are total of 21 steps in the installation wizard. During the installation progress, we selected "Server class" and "Single instance database installation" with the "Enterprise Edition" of Oracle Databases. We also gave "ORA12C" as my database Oracle System Identifier (SID). We specified the "Unicode" as the character set for the installation for enabling storage of multiple language groups. The database was installed in the directory "/ora01/app/oracle/ product/12.1.0/db 1/." The Oracle Enterprise Manager Database Express 12c the database installed by typing can access that we "https://fedora.dbaora.com:5500/em" on the web browser.

These steps were done for one database installation on the Fedora operating system.

B. SETTING UP AND MANAGING ORACLE DATA GUARD USING DATA GUARD COMMAND LINE INTERFACE

In this appendix, we describe how we set up and managed the Oracle Data Guard configurations using DGMGRL (Data Guard Command Line Interface). We used as a very well designed "hands on" lab, Data Guard Hands On Lab by Carpenter, which is included in the references section, in order to build up our system.

As described by [36], the setup and managing process can be expressed in nine main steps, which are preparing the primary database, creating the physical database, configuring data guard broker, changing transport mode using broker properties, changing protection mode to maximum availability, performing switchover from the primary database to the standby database, enabling the flashback database, performing manual failover from the primary database to the standby database, and enabling and using fast-start failover [36].

In the remainder of this appendix, we present how we set up our system and how we managed the Data Guard capability. Screenshots from the system are provided to illustrate each step.

1. Preparing the Primary Database

Figures 43 through 46 are given for preparing the primary database step.

Figure 43. Enabling Archiving and Force Logging

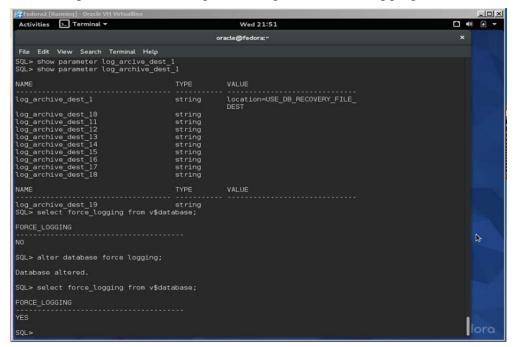


Figure 44. Adding 50MB Standby Redo Log Files

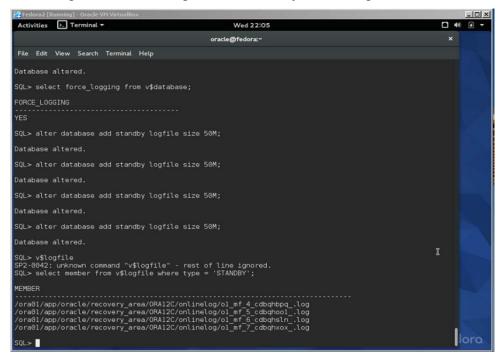


Figure 45. Tnsnames.ora File

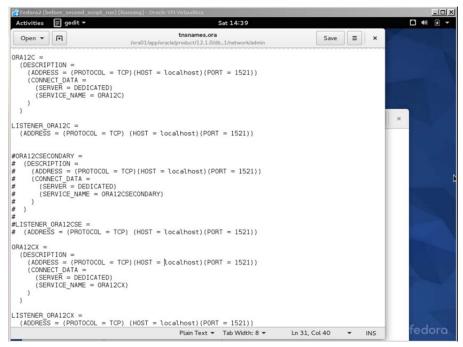
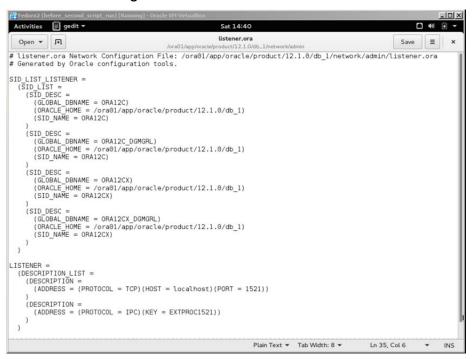


Figure 46. Listener.ora File



2. Creating the Physical Standby Database

Figures 47 and Figure 48 are given for preparing the primary database step.



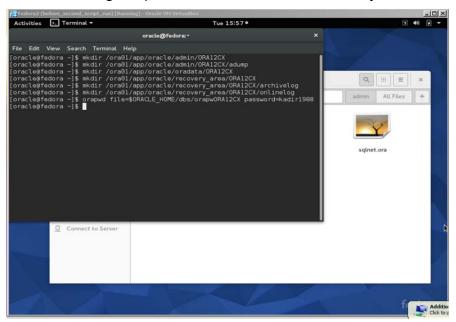
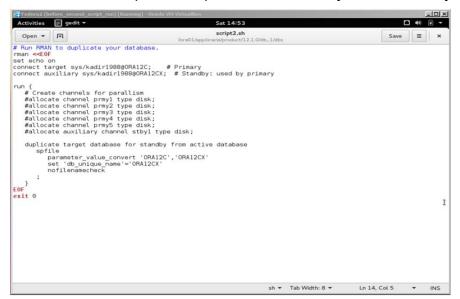


Figure 48. RMAN Script that Duplicates the Primary to the Standby



3. Configuring Data Guard Broker

Figures 49 through 59 are given for configuring data guard broker step.

Figure 49. Showing the Primary Database Broker Parameters

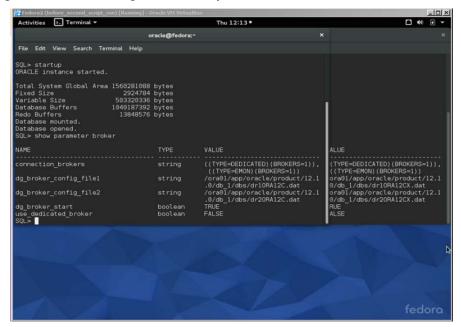


Figure 50. Showing the Standby Database Broker Parameters

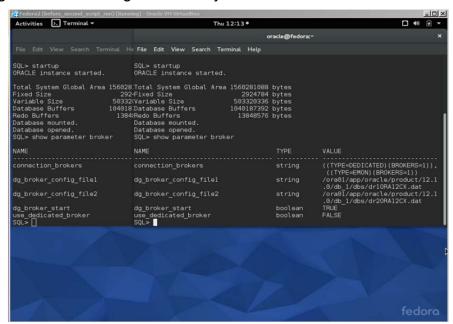


Figure 51. Creating the Broker Configurations-I

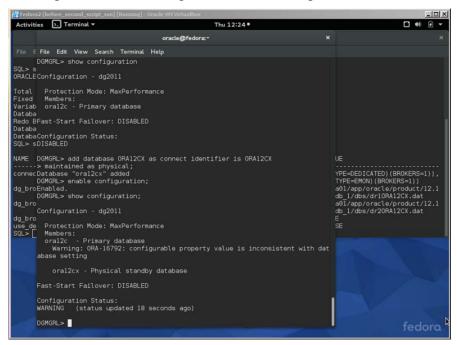


Figure 52. Creating the Broker Configurations-II

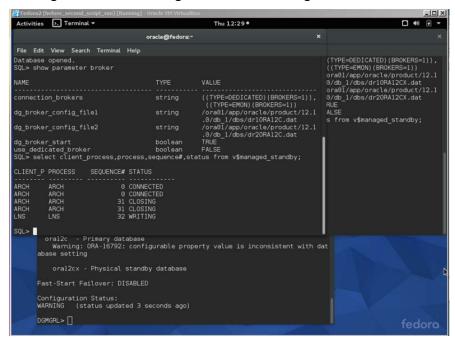


Figure 53. Creating the Broker Configurations-III

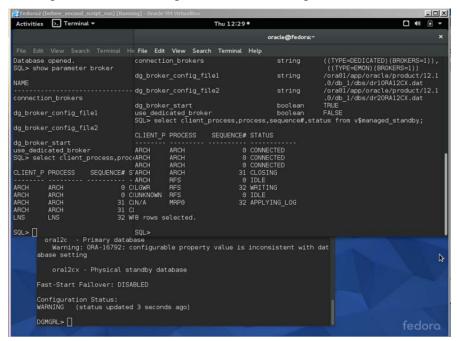


Figure 54. Showing the Primary Database Properties-I

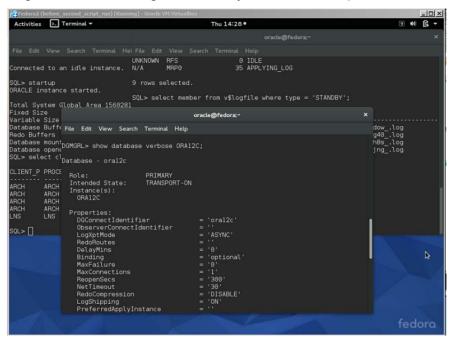


Figure 55. Showing the Primary Database Properties-II

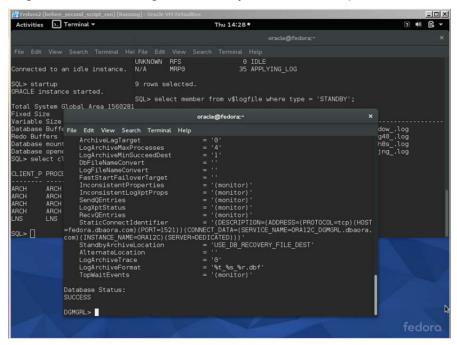


Figure 56. Showing the Standby Database Properties-I

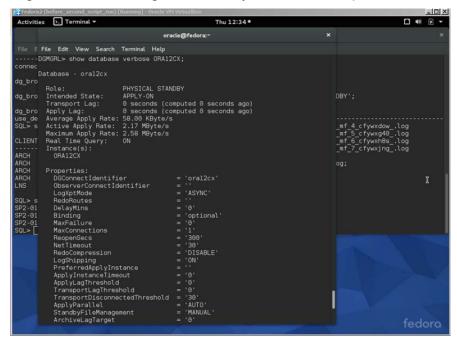


Figure 57. Showing the Standby Database Properties-II

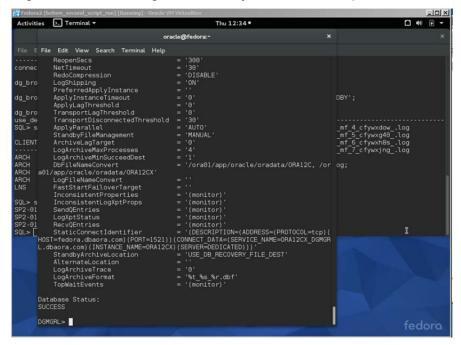
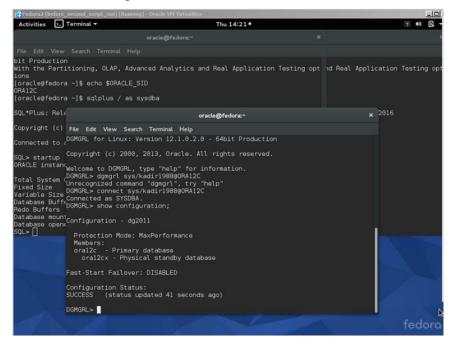


Figure 58. The Configuration of Data Guard, with MaxPerformance



Activities Terminal True 00:53 True | Tru

Figure 59. The Configuration of Data Guard, with MaxAvailability

4. Changing Transport Mode Using Broker Properties

Members: oral2c - Primary database oral2cx - Physical standby database

Configuration Status: SUCCESS (status updated 34 seconds ago)

Fast-Start Failover: DISABLED

DGMGRL>

Figures 60 through 65 are given for changing transport mode using broker properties step.

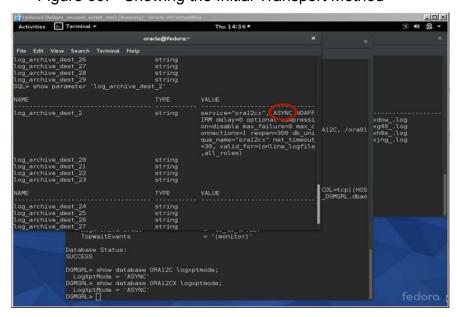


Figure 60. Showing the Initial Transport Method

Figure 61. Modifying the Transport Mode to Synchronous

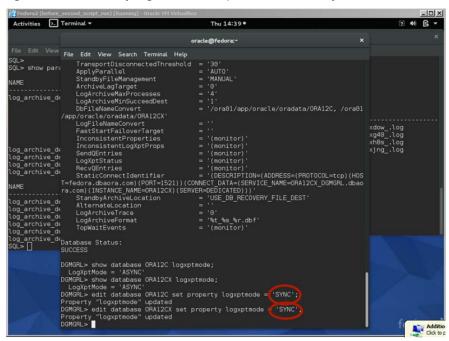


Figure 62. Verifying the Transport Method

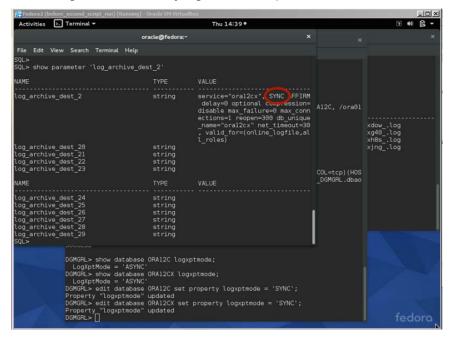


Figure 63. Creating a Redo Gap Between Primary and Standby

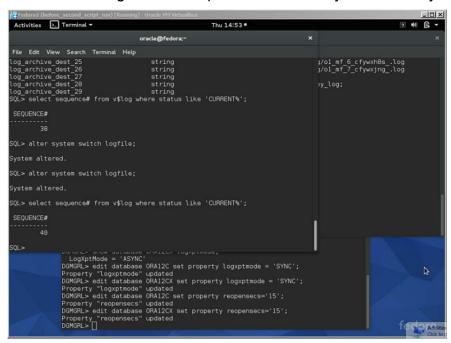


Figure 64. Restarting the Standby Database

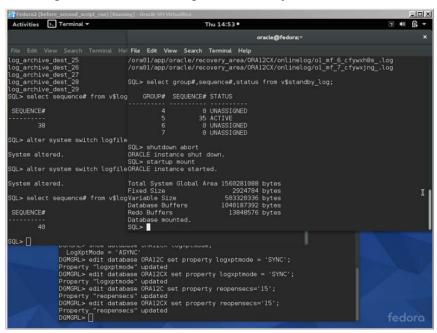
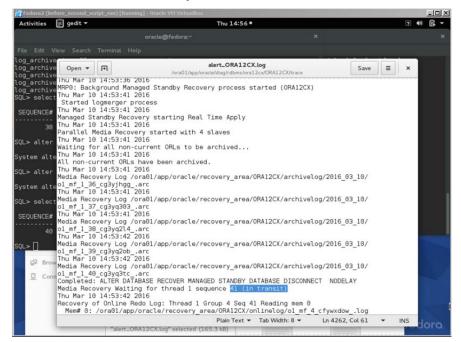


Figure 65. Verifying Primary and the Standby Databases Automatic Synchronization



5. Changing Protection Mode to Maximum Availability

Figures 66 and Figure 67 are given for changing protection mode to maximum availability step.



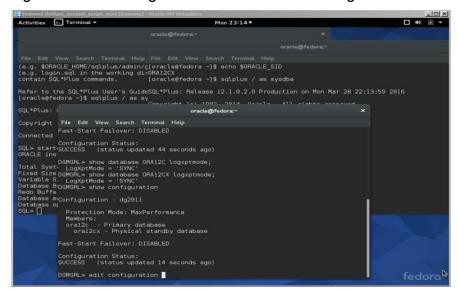
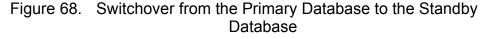


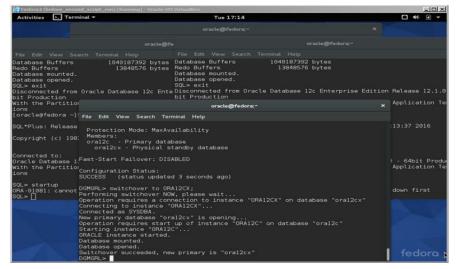
Figure 67. Showing the Protection Mode Configuration-II

6. Performing Switchover from the Primary Database to the Standby Database

DGMGRL>

Figure 68 is given for performing switchover from the primary database to the standby database step.

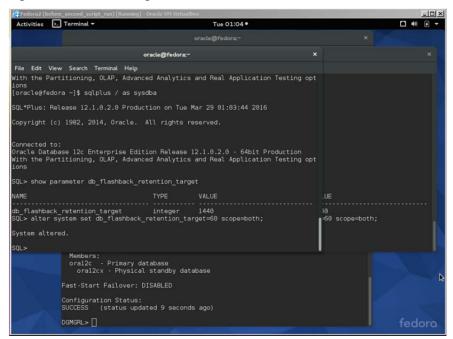




7. Enabling Flashback Database

Figures 69 through 72 are given for enabling flashback database step.

Figure 69. Showing Parameters for Flashback Database



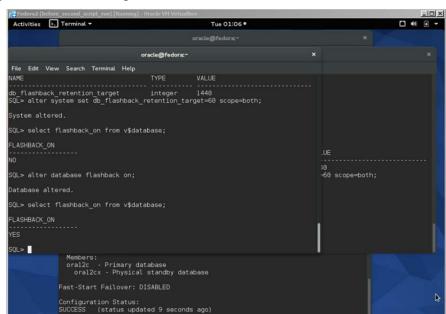
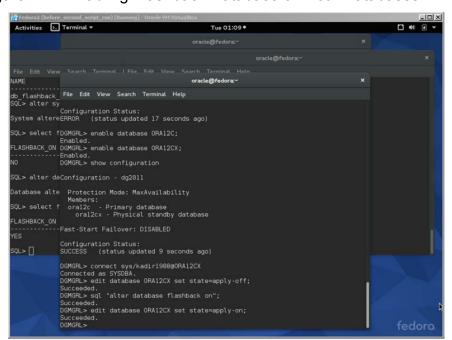


Figure 70. Enabling Flashback Database on Both Databases-I

Figure 71. Enabling Flashback Database on Both Databases-II



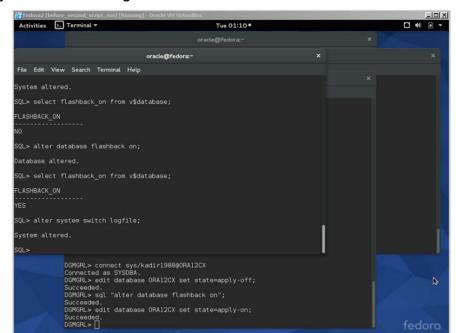


Figure 72. Enabling Flashback Database on Both Databases-III

8. Performing Manual Failover from the Primary Database to the Standby Database

Figures 73 through 79 are given for performing manual failover from the primary database to the standby database step.

Figure 73. Shutting Down the Primary and Showing Configurations

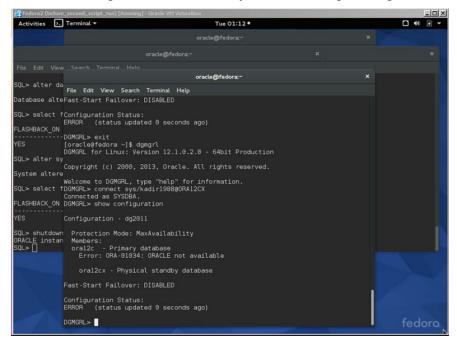


Figure 74. Successful Manual Failover

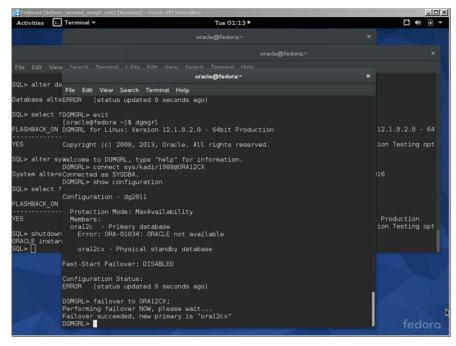


Figure 75. Verifying that the Standby Database Became the Primary

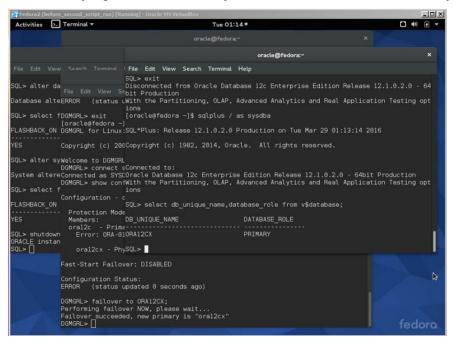


Figure 76. Starting Up the Standby Database (Previously the Primary)

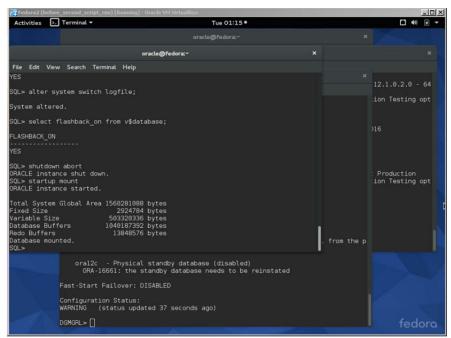


Figure 77. The Broker Initiates Reinstatement of Database that Shut Down Unexpectedly

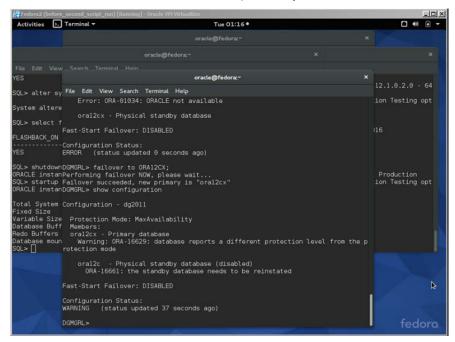


Figure 78. Success after Reinstating the Database that Shut Down Unexpectedly

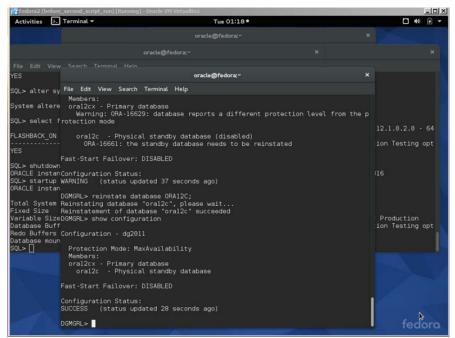
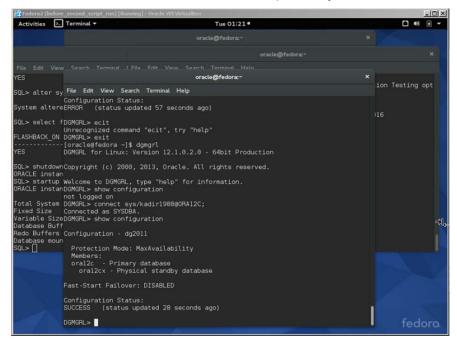


Figure 79. Success in Configuration after Switchover to the Database that Shut Down Unexpectedly



9. Enabling and Using Fast-Start Failover

Figures 80 through 89 are given for enabling and using fast-start failover step.

Figure 80. Changing the Fast-Start Failover Threshold

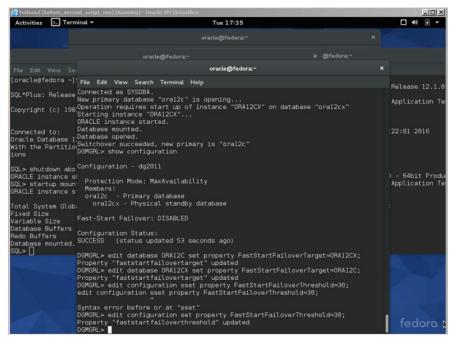
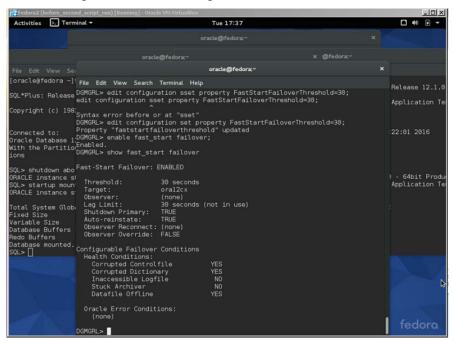


Figure 81. Enabling Fast-Start Failover



Activities Terminal Toracle@fedorx"

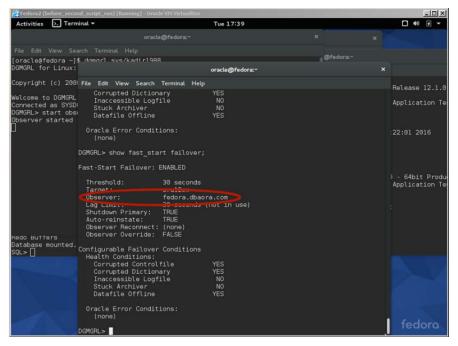
Activities Terminal Toracle@fedorx"

File Edit View Search Terminal Help

| Convert | Con

Figure 82. Starting the Observer

Figure 83. Verifying the Observer Is Started



Activities Terminal Terminal Terminal Terminal Help

bit Production
With the Partitioning, OLAP, Advanced Analytics and Real Application Testing opt
ions
[oracle@fedorac] \$ sqlplus / as sysdba

SQL*Plus: Release 12.1.9.2.9 Production on Tue Mar 29 17:49:41 2016

Copyright (c) 1982, 2014, Oracle. All rights reserved.

Connected to:

Oracle Database 12c Enterprise Edition Release 12.1.9.2.9 - 64bit Production
With the Partitioning, OLAP, Advanced Analytics and Real Application Testing opt
ions

SQL*Plus: Release 12.1.9.2.9 Production on Tue Mar 29 17:40:41 2016

Copyright (c) 1982, 2014, Oracle. All rights reserved.

Connected to:
Oracle Database 12c Enterprise Edition Release 12.1.9.2.9 - 64bit Production
With the Partitioning, OLAP, Advanced Analytics and Real Application Testing opt
ions

SQL* select db_unique_name, database_role from v\$database;

DB_UNIQUE_NAME

DATABASE_ROLE

COPYRIDED Control Tile
COPY

Figure 84. Shutting Down the Primary Database for Fast-Start Failover

Figure 85. Examining the Actions in the Observer during the Fast-Start Failover

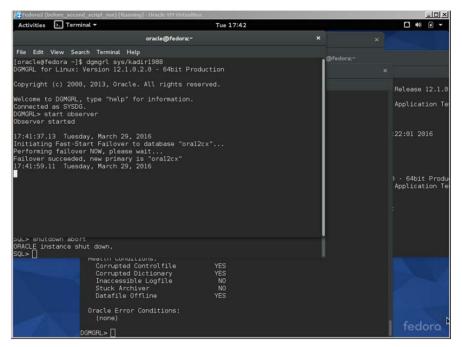


Figure 86. The Configuration after Fast-Start Failover

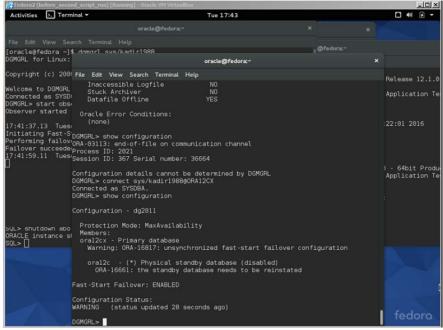


Figure 87. Restarting the Previous Primary Database

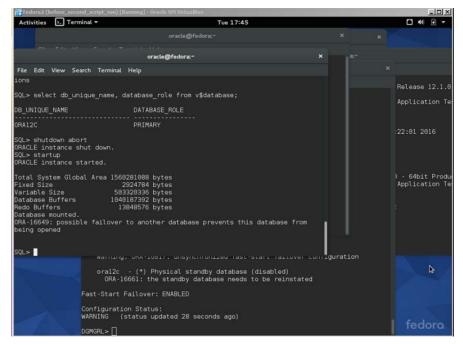


Figure 88. Examining the Actions in the Observer after Restarting the Previous Primary Database

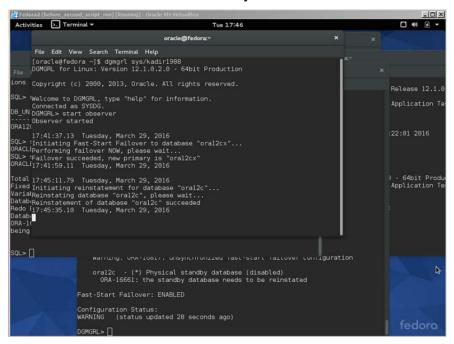
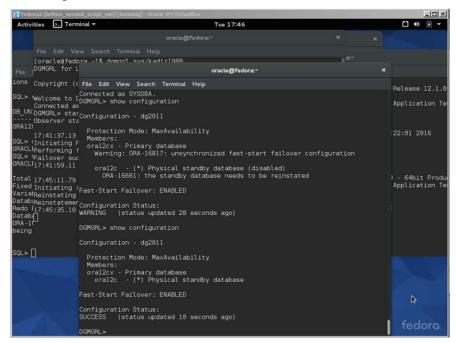


Figure 89. Success after Fast-Start Failovers



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